



Experiences on K computer from a topic focused on the large-scale eigenvalue solver project

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Joint work with

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24th Workshop for Sustained Simulation Performance,
Stuttgart HLRS Aquarium, 5-6Dec, 2016



1. Quick Overview of Project

- Past and Present of EigenExa
- Diagonalization of a 1million x 1million matrix on K computer

2. The latest updates

3. Future direction

4. Summary

Key topic is

How to remove three walls;

i) Memory bandwidth, ii) Network Latency, iii) Parallelism.

H4ES (2011-2016)

1. Eigenvalue problem is of significance in many scientific and engineering fields

$$Ax = \lambda Bx$$

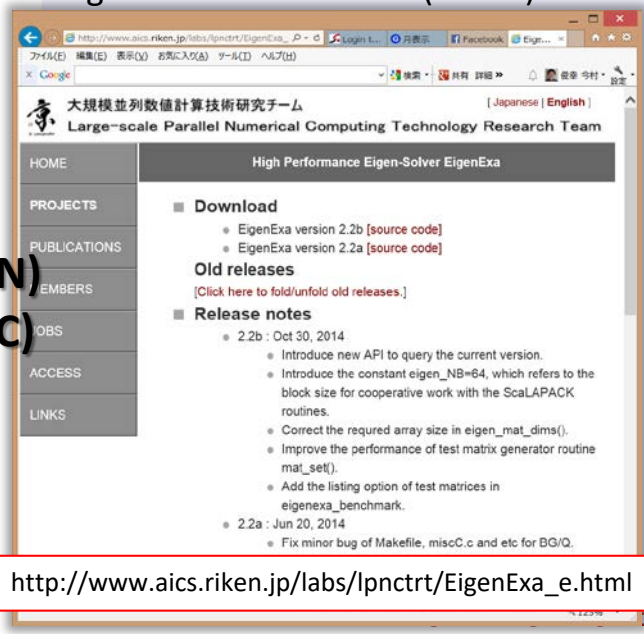
2. In practical simulation, Sparse and dense solver must be tightly cooperating

3. Currently, collaboration with ESSEX under the joint initiative of DFG-JST-ANR (2016-2018).

Prof. Sakurai Team 'H4ES'

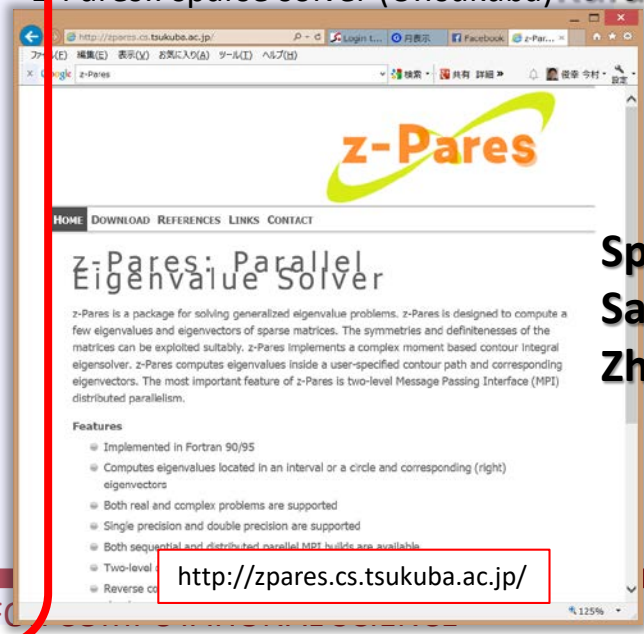
Application:
Hoshi(Tottori)

EigenExa:: dense solver (RIKEN)



http://www.aics.riken.jp/labs/lpnctr/EigenExa_e.html

z-Pares:: sparse solver (U.Tsukuba) Kuramashi(Tsukuba)



<http://zpare.cs.tsukuba.ac.jp/>

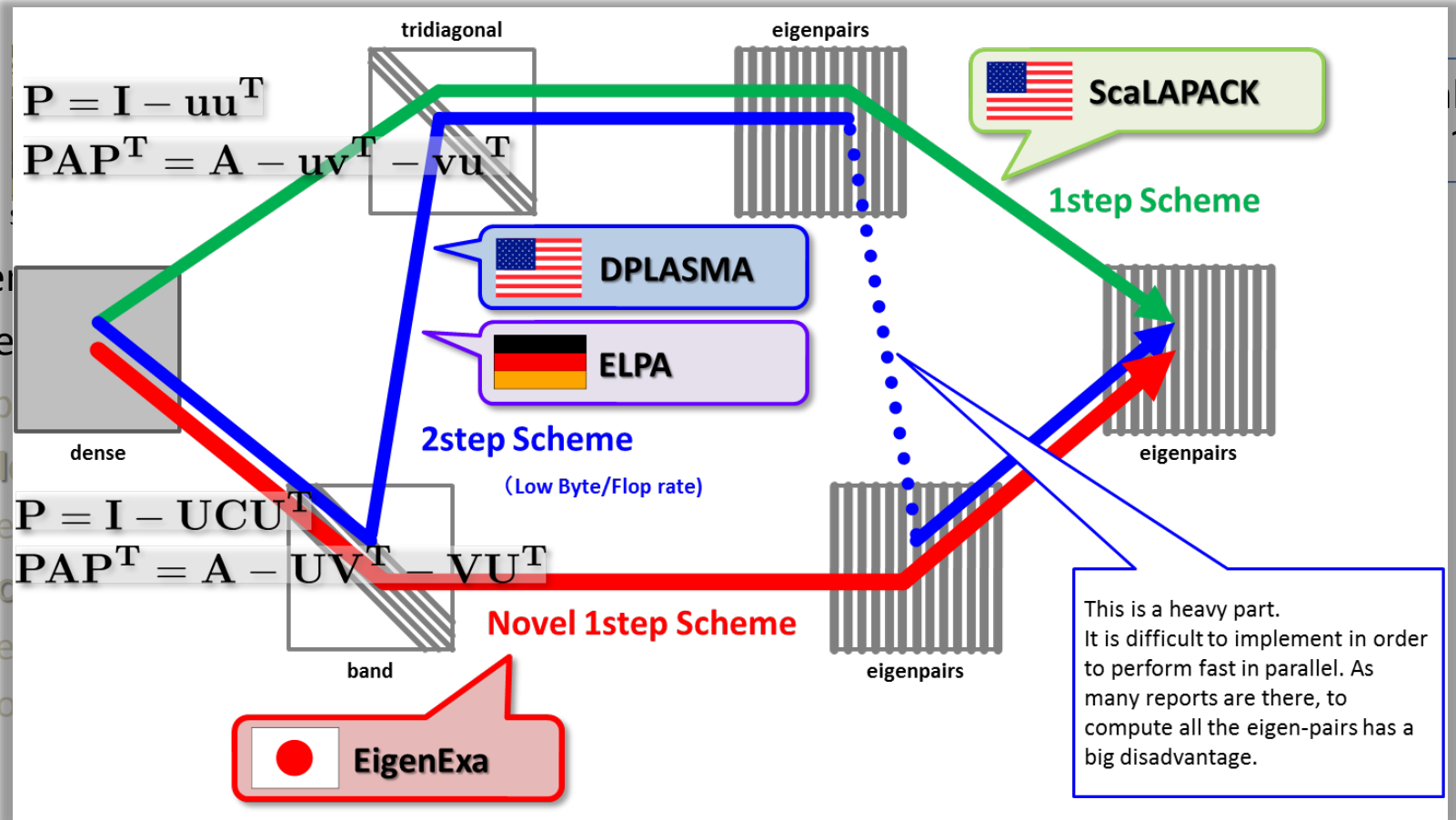
Dense:
Imamura(RIKEN)
Yamamoto(UEC)

Sparse & LS:
Sakurai(Tsukuba)
Zhang(Nagoya)

EigenExa Project

- Project itself is old...
 - Earth Simulator version was published in SC06. and the speakers continue to update it approximately 10 years. Partly funded by another CREST organized by Prof. Yagawa.

- Current
- Eigen
- Two b
- 1. Bl
- Re
- 2. Co
- Re
- Co



le
??

```

k_1 = i - i_base
k_2 = m0

L = i-1
n = eigen_translate_g2l(L, x_nnod, x_inod)

!$OMP MASTER
  call eigen_vector_zeropad_x( v_x(1), L )
!$OMP END MASTER

!$OMP BARRIER
  prod_uv = u_t(1)
  if ( k_2 <= k_1 ) then
    if ( beta /= ZERO ) then
      alpha = prod_uv/(2*beta)
!$OMP DO
  do j_1=1,n
    v_x(j_1) = (v_x(j_1)-alpha*u_x(j_1))/beta
  end do      ! j_1
!$OMP ENDDO
  end if
  else
!
! v=v-(UV+VU)u
!
  l_4 = MOD(k_2-k_1, 3)+k_1+1
  LX = 64      ! L1_LSIZE*L1_WAY/16

  LL = (n-1)/y_nnod+1
  LL = ((LL-1)/2+1)*2

  alpha = ddot( k_2-k_1, u_t(2), 2, u_t(3), 2 )
  prod_uv = prod_uv - 2*alpha

  if ( n > VTOL ) then
    jj_2 = 1+LL*(y_inod-1)
    jj_3 = MIN(n, LL*y_inod)
  else
    jj_2 = 1
    jj_3 = n
  endif

!$OMP DO
  do jj_1=jj_2,jj_3,LX
    j_2 = jj_1; j_3 = MIN(jj_1+LX-1, jj_3)
    if(l_4-1==k_1+1)then
      l_1 = k_1+1    ! 0
      j = l_1-k_1

      u0 = u_t(2*(j+0)-1+1)

```

```

v0 = u_t(2*(j+0)-0+1)
do j_1=j_2,j_3
  w0 = v_x(j_1)
  ux0 = ux(j_1,l_1+0)
  vx0 = vx(j_1,l_1+0)
  w0 = w0-ux0*u0-vx0*v0
  v_x(j_1) = w0
end do      ! j_1
end if
if(l_4-2==k_1+1)then
  l_1 = k_1+1    ! 1

  j = l_1-k_1

  u0 = u_t(2*(j+0)-1+1)
  v0 = u_t(2*(j+0)-0+1)
  u1 = u_t(2*(j+1)-1+1)
  v1 = u_t(2*(j+1)-0+1)
  do j_1=j_2,j_3
    w0 = v_x(j_1)
    ux0 = ux(j_1,l_1+0)
    vx0 = vx(j_1,l_1+0)
    w0 = w0-ux0*u0-vx0*v0
    ux1 = ux(j_1,l_1+1)
    vx1 = vx(j_1,l_1+1)
    w0 = w0-ux1*u1-vx1*v1
    v_x(j_1) = w0
  end do      ! j_1
end if
do l_1=l_4,k_2,3    ! 2

  j = l_1-k_1

  u0 = u_t(2*(j+0)-1+1)
  v0 = u_t(2*(j+0)-0+1)
  u1 = u_t(2*(j+1)-1+1)
  v1 = u_t(2*(j+1)-0+1)
  u2 = u_t(2*(j+2)-1+1)
  v2 = u_t(2*(j+2)-0+1)
  do j_1=j_2,j_3
    w0 = v_x(j_1)
    ux0 = ux(j_1,l_1+0)
    vx0 = vx(j_1,l_1+0)
    w0 = w0-ux0*u0-vx0*v0
    ux1 = ux(j_1,l_1+1)
    vx1 = vx(j_1,l_1+1)
    w0 = w0-ux1*u1-vx1*v1
    ux2 = ux(j_1,l_1+2)
    vx2 = vx(j_1,l_1+2)
    w0 = w0-ux2*u2-vx2*v2

```

```

  v_x(j_1) = w0
  end do      ! j_1
  end do      ! l_1
!$OMP ENDDO

  if ( beta /= ZERO ) then
    alpha = prod_uv/(2*beta)
!$OMP DO
  do j_1=jj_2,jj_3
    v_x(j_1) = (v_x(j_1)-alpha*u_x(j_1))/beta
  end do      ! j_1
!$OMP ENDDO
  end if
!$OMP MASTER
  if ( n > VTOL ) then
    call allgather_db1(v_x(jj_2), v_t, LL, 1, y_COMM_WORLD)
    j_3 = eigen_loop_end(L, x_nnod, x_inod)
    v_x(1:j_3) = v_t(1:j_3)
  end if
!$OMP END MASTER
  end if
!$OMP BARRIER

!$OMP MASTER
  v_n = (v_n-alpha*u_n)/beta
  x_owner_nod = eigen_owner_node (L, x_nnod, x_inod)
  x_pos = eigen_translate_g2l(L, x_nnod, x_inod)
  if ( x_inod == x_owner_nod ) then
    v_x(x_pos) = v_n
  end if
  if ( kk == 0 ) then
    call datacast_db1( v_y(1), v_x(1), u_t(1), v_t(1), x_pos, 2 )
  end if
  call eigen_vector_zeropad_x( v_x(1), L )
  if ( kk == 0 ) then
    call eigen_vector_zeropad_y( v_y(1), L )
  end if
!$OMP END MASTER

```

Simulation codes

- 8 Applications
 - **Platypus QM/MM**: gives the precise analysis for a biological polymer such as a kinase reaction mechanism by introducing the electron state effect to the molecular mechanics (MM) approach
 - **RSDFT**: is an ab-initio program with the real-space difference method and a pseudo-potential method
 - **PHASE**: is a Computer Software for Band Calculations based on First-principles Pseudo-potential Method
 - **ELSES**: is large-scale atomistic simulation with quantum mechanical freedom of electrons manipulating a large Hamiltonian matrix.
 - **NTChem**: is a high-performance software package for the molecular electronic structure calculation for general purpose on the K computer
 - **Rokko**: Integrated Interface for libraries of eigenvalue decomposition
 - **LETKF**: data assimilation for atmospheric and oceanic systems
 - **POD**: proper orthogonal decomposition (POD) to compress data for example video data

Material Science 1

- Hasegawa, Y, etc. First-principles calculations of electron states of a silicon nanowire with 100,000 atoms on the K computer, SC11, *Gordon Bell Prize Winner*, 2011.
- First principle electronic structure simulation
- Implemented real space method (not use FFT)
- Main procedure are conjugate-gradient method, ortho-normalization, and subspace diagonalization
- Use MPI + OpenMP
- Parallelized in grid space and orbitals
- Use DGEMM, and eigensolver library (ScaLAPACK, EIGEN)

Silicon nanowire 39,696 atoms
10648-atom cell of Si crystal and its electron density

RIKEN ADVANCED INSTITUTE FOR COMPUTATIONAL SCIENCE

Material Science 2

Imachi and Hoshi (U.Tottori, JST-CREST)

- **ELSES** (<http://www.elses.jp/>): Extra-large-scale electronic-structure simulation code
- (Nano polycrystalline diamond)
- ⇒ eigenstate obtained by solving a 430K dimensional matrix

17nm

(*) ナノ多結晶ダイヤモンド (超強度材料)
[1] (合成@愛媛大)
T. Irifune, et al., Nature 421, 599 (2003)
[2] (理論研究) T. Hoshi, et al., J. Phys. Soc. Jpn. 82, 023710 (2013)
[3] 製品化 (住友電工, 2012)

50nm

Big Data Science

- World record scale global ensemble data assimilation by LETKF
- Performance: 263 TFLOPS (>44% theoretical peak) with taking advantage of 4608 nodes of the K computer (590 TFLOPS)
- 'Using the efficient eigenvalue solver for the K computer, the LETKF computations are accelerated by a factor of 8, allowing a 3 week experiment of 10,240-member LETKF with an intermediate AGCM for the first time.'

T. Miyoshi, K. Kondo and T. Yamamura, The 10,240-member ensemble Kalman filtering with an intermediate AGCM, Geophysical Research Letters, Vol 41, 14, pp.5264-5271 (2014) DOI: 10.1002/2014GL060863

10240 members w/o localization

100 members 10240 members

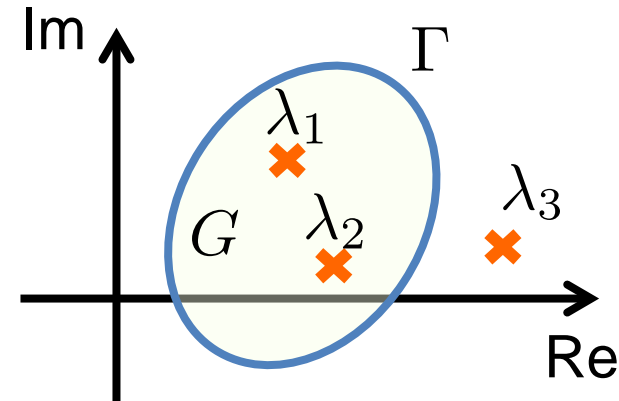
アンサンブルデータ同化による水蒸気量の誤差を表すヒストグラム
真値はガス分布 (正準分布) 関数、10,240個のアンサンブル (右) により、100個のアンサンブル (左) では確認できなかった大気状態のほとんどの非ガウス性を初めて直接確認。

アンサンブルデータ同化による水蒸気量の相関マップ
黄色い線の境界内の水蒸気量の観測データに対する各地点の水蒸気量の相関係数を示す。アンサンブル員数が増えることで観測ノイズが抑えられ、観測データが1万km平方に及ぼす影響まで確認できた。

Sakurai-Sugiura eigenvalue solver

- Contour integral for a rational function

$$\frac{1}{2\pi i} \oint_{\Gamma} \sum_{i=1}^n \frac{\nu_i}{z - \lambda_i} dz = \sum_{\lambda_i \in G} \nu_i$$



- Spectral decomposition of $(zB - A)^{-1}B$:

$$(zB - A)^{-1}B = \sum_{i=1}^n \frac{P_i}{z - \lambda_i}$$

λ_i : eigenvalue, P_i : spectral projection with respect to λ_i
(for simplicity, we consider the case that λ_i is simple)

Localization of spectral decomposition using contour integral

$$P_{\Gamma} = \frac{1}{2\pi i} \oint_{\Gamma} (zB - A)^{-1}B dz = \sum_{\lambda_i \in G} P_i$$

z-Pares

- Implemented in Fortran 95 and MPI
 - C interface will be provided
- Provides subroutines for
 - A,B real symm, B positive definite
 - A,B Hermitian, B positive definite
 - A,B real unsymmetric
 - A,B complex non-Hermitian
- Provides efficient implementation for standard problem
- Dependencies
 - BLAS/LAPACK
 - MUMPS* (Optional)

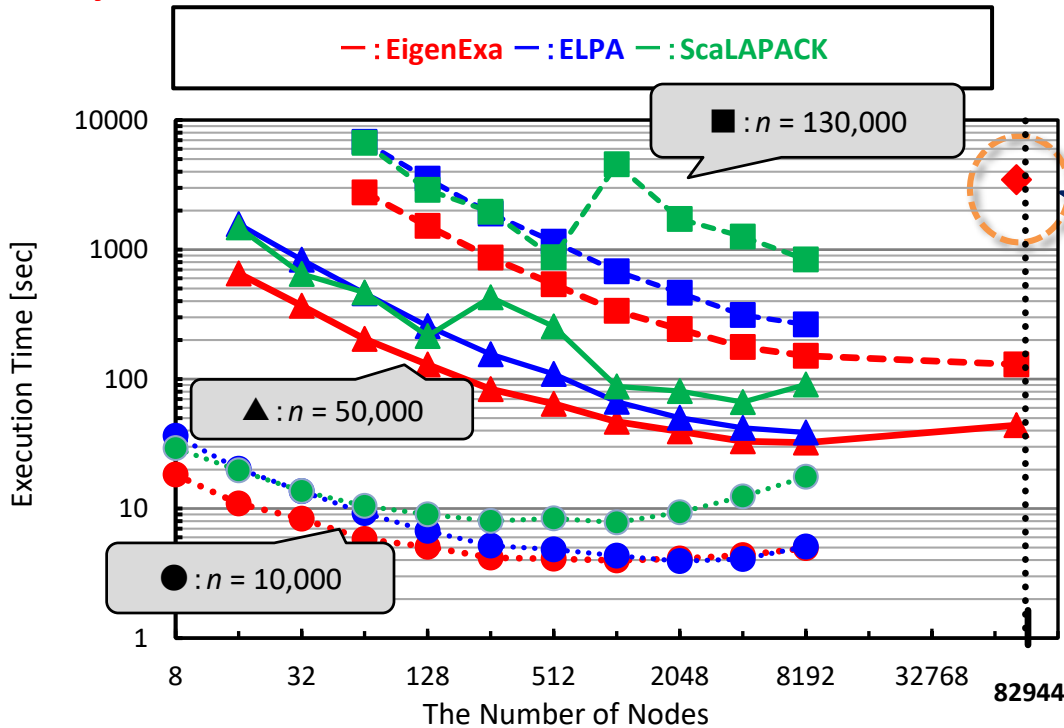


$$Ax = \lambda Bx$$

* MPI distributed parallel sparse direct linear solver

World Largest Dense Eigenvalue Computation

- We have successfully done a world largest-scale dense eigenvalue benchmark (one million dimension) by EigenExa taking advantage of the overall nodes (82,944 processors) of K computer in 3,464 seconds. Our EigenExa achieves 1.7 PFLOPS (16% of the K computer's peak performance).
- Feasibility and Reliability for algorithm and library are confirmed, especially assumed on a post-K system.**



: $n = 1,000,000$
 EigenExa solves a world largest-scale problem.
 (1.7 PFLOPS, 16% of K computer's theoretical peak performance)

$$\max_i \|Av_i - \lambda_i v_i\|_2 / \|A\|_F = 3.1 \times 10^{-13}$$

$$\|V^T V - I\|_F = 2.1 \times 10^{-10}$$

- ✓ n is the dimension of problems.
- ✓ 1 MPI process * 8 threads per node.
- ✓ Test matrices are randomly generated.

Specification of K computer

- Peak performance: 10.6 PFLOPS
- Num. of Nodes: 82,944
- Performance/node: 128 GFLOPS (One octa-core SPARC 64 VIIIfx)
- Network: Tofu interconnect (6D mesh-torus)

Related performance report is T.Fukaya, TI. "Performance evaluation of the EigenExa eigensolver on Oakleaf-FX: tridiagonalization versus pentadiagonalization", PDSEC2015

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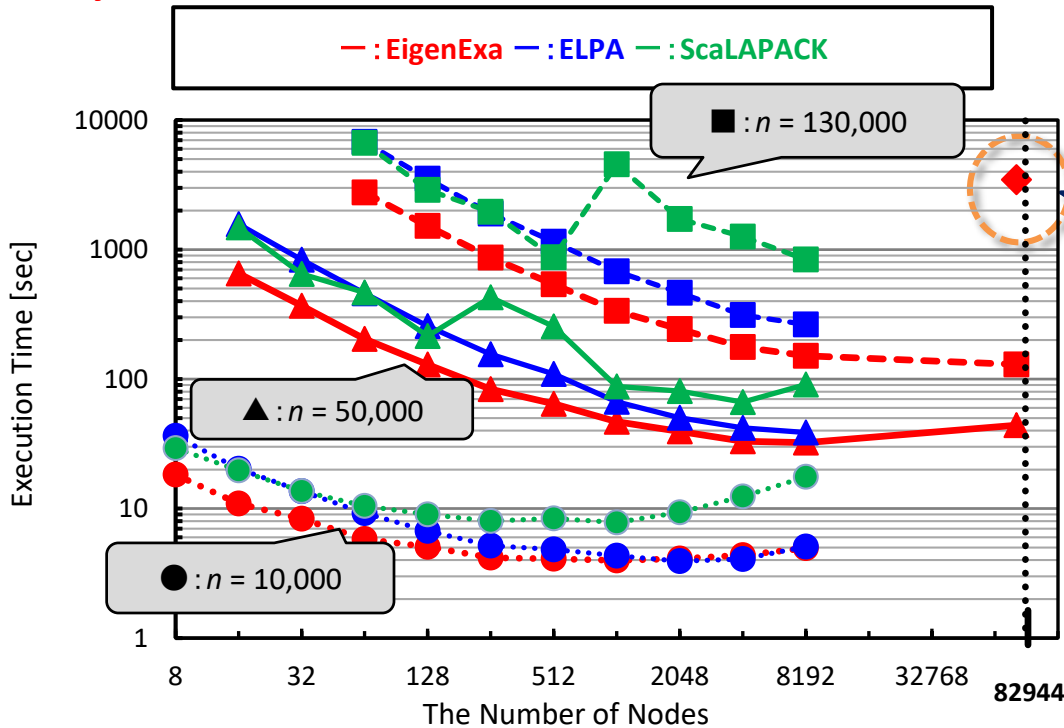
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What we got from the ultra-scale experiments

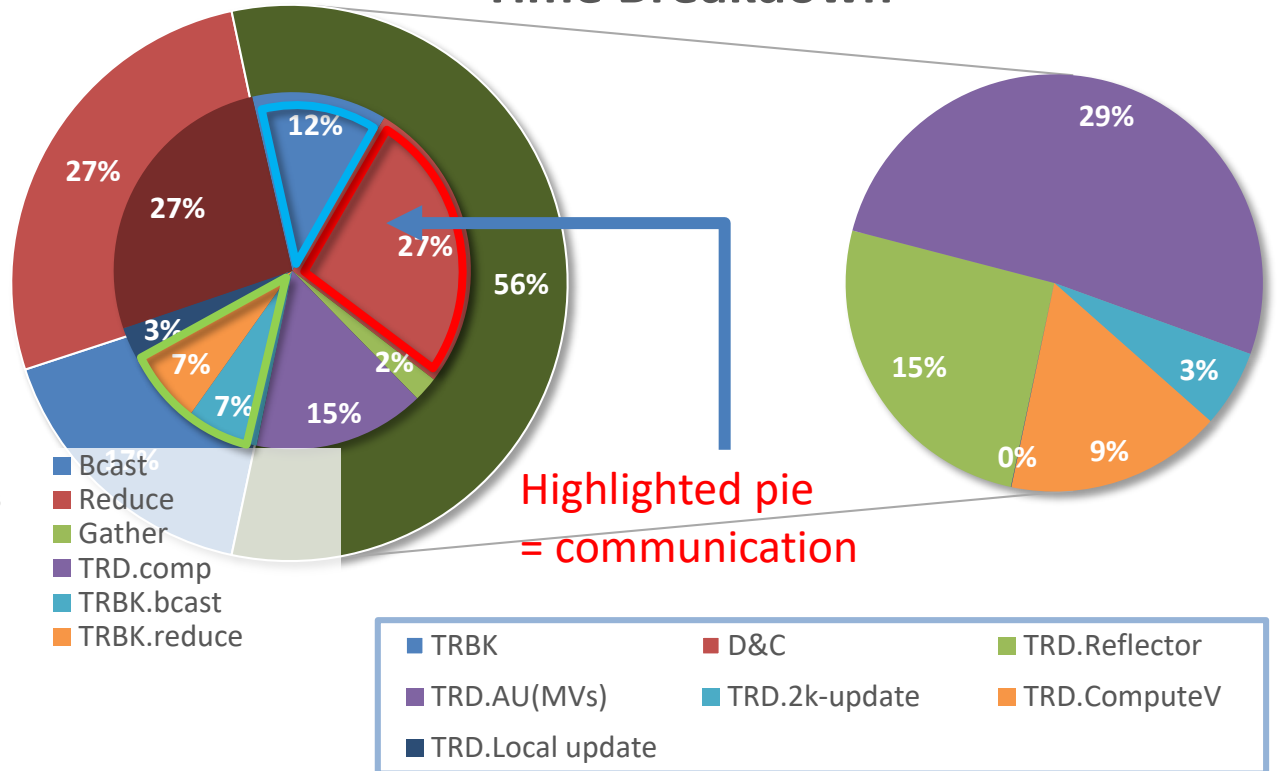
```

NUM.OF.PROCESS= 82944 ( 288 288 )
NUM.OF.THREADS= 8
calc (u,beta) 503.0970594882965
mat-vec (Au) 1007.285000801086 661845.1244051798
2update (A-uv-vu) 117.4089198112488
5678160.294281102
calc v 0.0000000000000000
v=v-(UV+VU)u 328.3385872840881
UV post reduction 0.6406571865081787
COMM_STAT
BCAST :: 424.3022489547729
REDUCE :: 928.1299135684967
REDIST :: 0.0000000000000000
GATHER :: 78.28400993347168
TRD-BLK 1000000 1968.435860157013
677356.7583893638 GFLOPS
TRD-BLK-INFO 1000000 48
before PDSTEDC 0.1448299884796143
PDSTEDC 905.2210271358490
MY-REDIST1 1.544256925582886
MY-REDIST2 14.75343394279480
RERE1 4.861211776733398E-02
COMM_STAT
BCAST :: 4.860305786132812E-02
REDUCE :: 2.155399322509766E-02
REDIST :: 0.0000000000000000
GATHER :: 0.0000000000000000
PDGEMM 532.6731402873993 5417097.565200153
GFLOPS
D&C 921.8044028282166 3130319.580211733 GFLOPS
TRBAK= 573.9026420116425
COMM= 533.7601048946381
573.9026420116425 3484911.644577213 GFLOPS
182.3303561210632 5484550.248648792 GFLOPS
152.0370917320251 6577342.335399065 GFLOPS
0.1022961139678955 7.379654884338379
COMM_STAT
BCAST :: 229.3666801452637
REDUCE :: 234.4477448463440
REDIST :: 0.0000000000000000
GATHER :: 0.0000000000000000
TRBAKWY 573.9029450416565
TRDBAK 1000000 573.9216639995575 3484796.141101135
GFLOPS
Total 3464.162075996399 1795203.448396145 GFLOPS
Matrix dimension = 1000000
Internally required memory = 480502032 [Byte]
Elapsed time = 3464.187163788010 [sec]

```

World Largest Dense Eigenvalue Computation

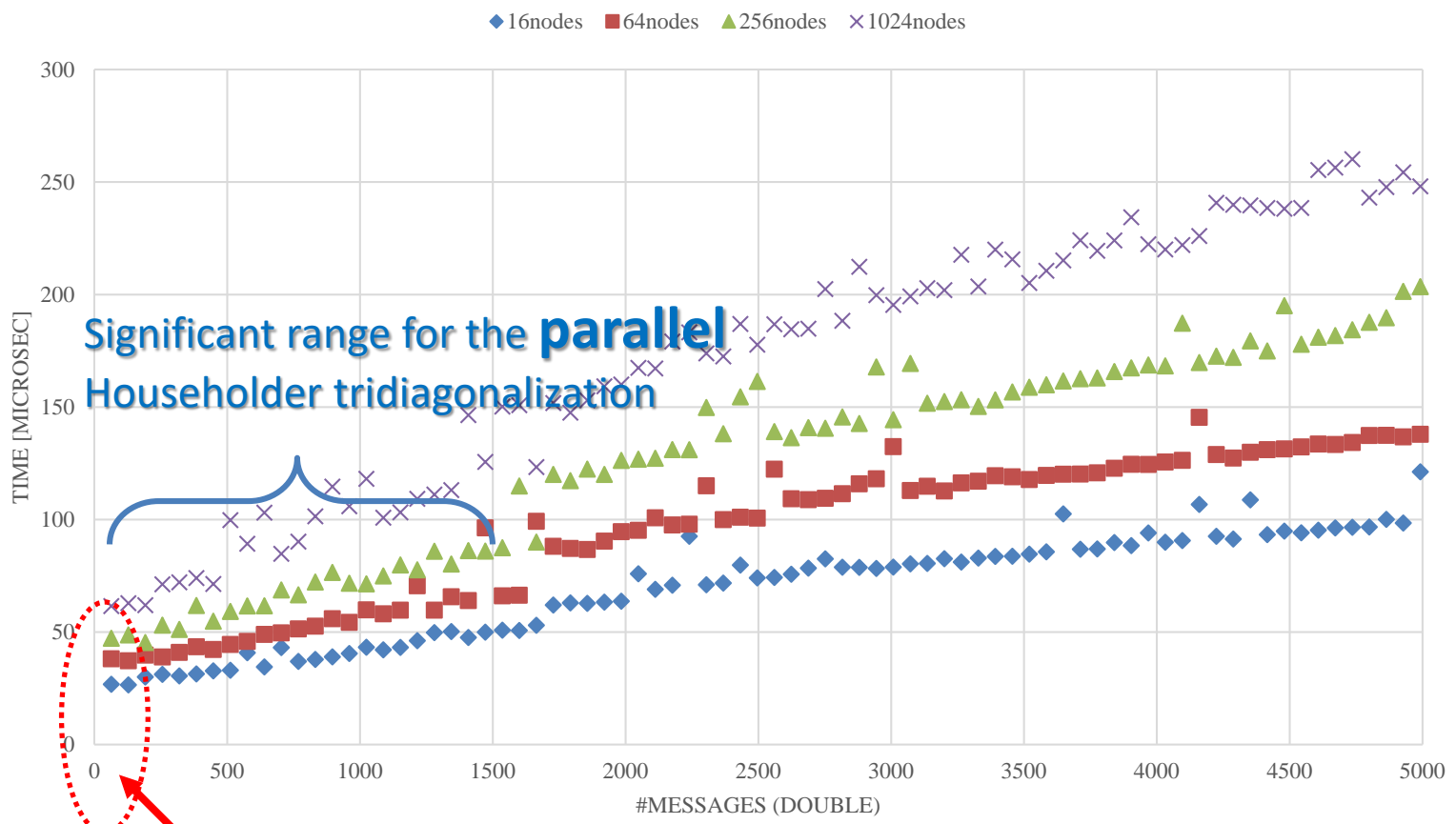
Time Breakdown



Related performance report is
H.Imachi, T. Hoshi "Hybrid Numerical Solvers for Massively Parallel Eigenvalue Computations and Their Benchmark with Electronic Structure Calculations", Journal of Information Processing Vol. 24 (2016) No. 1 pp. 164-172

Allreduce is an expensive operation

BENCHMARK OF MUTI-MPI_ALLREDUCE ON K COMPUTER



startup cost is 25~60 microseconds!

~ equivalent to pure time of allreduce with 1500 words

CA for EigenExa

- **Communication Avoiding algorithm**

- Blocking technique, increasing locality by data replication, and exchange the operation order.
- Introducing an extended form of vector 'A'.
- Computing Au and $u^T u$, simultaneously.

TI, etc, "CAHTR: Communication-Avoiding Householder Tridiagonalization", ParCo15

TI, "Parallel dense eigenvalue solver and SVD solver for post-petascale computing systems", PMAA16

Naïve version of the 2-sided HH trans.

$$s = \text{sign}(\|u\|, -(u, e))$$

$$u := u - se$$

$$v = Au$$

$$[C_U; C_V] = [U^T; V^T]u$$

$$v := v - (UC_V + VC_U)$$

$$f = (u, v)$$

$$v := v - afu$$

Single or several word allreduce

CA for EigenExa

Communication Avoiding algorithm

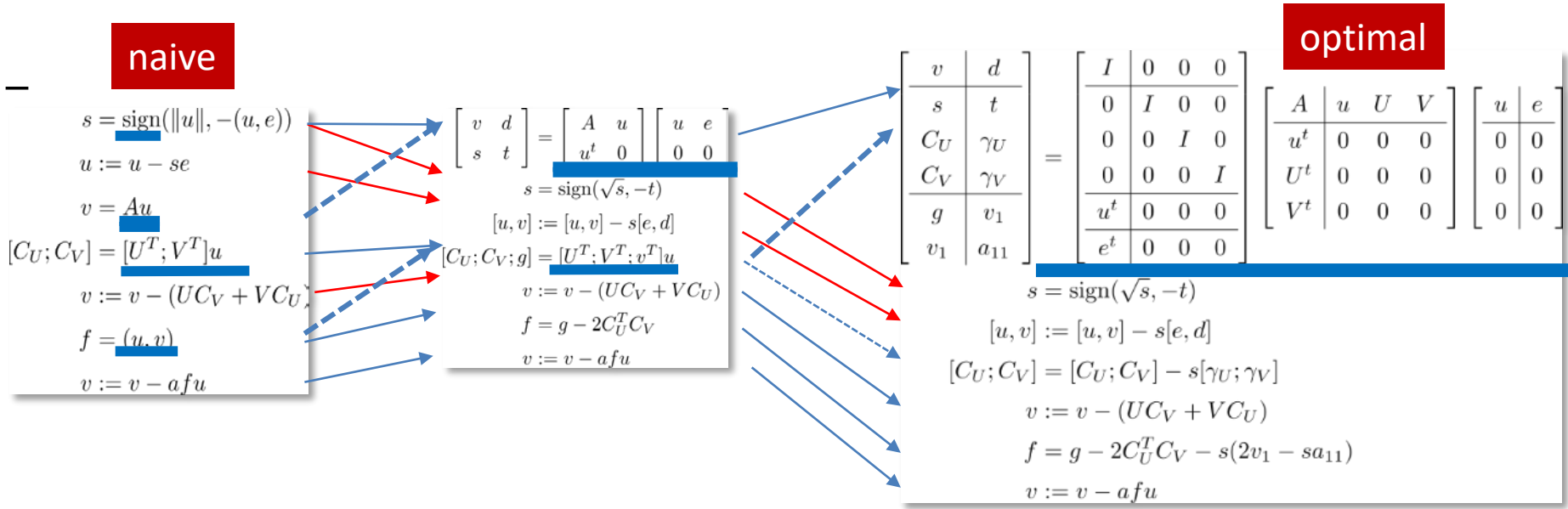
- Blocking technique, increasing locality by data replication, and exchange the operation order.
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TI, "Parallel dense eigenvalue solver and SVD solver for post-petascale computing systems", PMAA16

Principles :

Distributive Law & Exchange order & Introducing correction terms & Combine couples of collective operations into one



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Current status

- 1million x 1million
- Introduction of CA

Eigen-Exa

Eigen-Exa2 2020

2011



Implementation of K

2013



Performance
Evaluation

2016



Porting to other peta-systems

We are here

Post-K (10P>) supercomputer



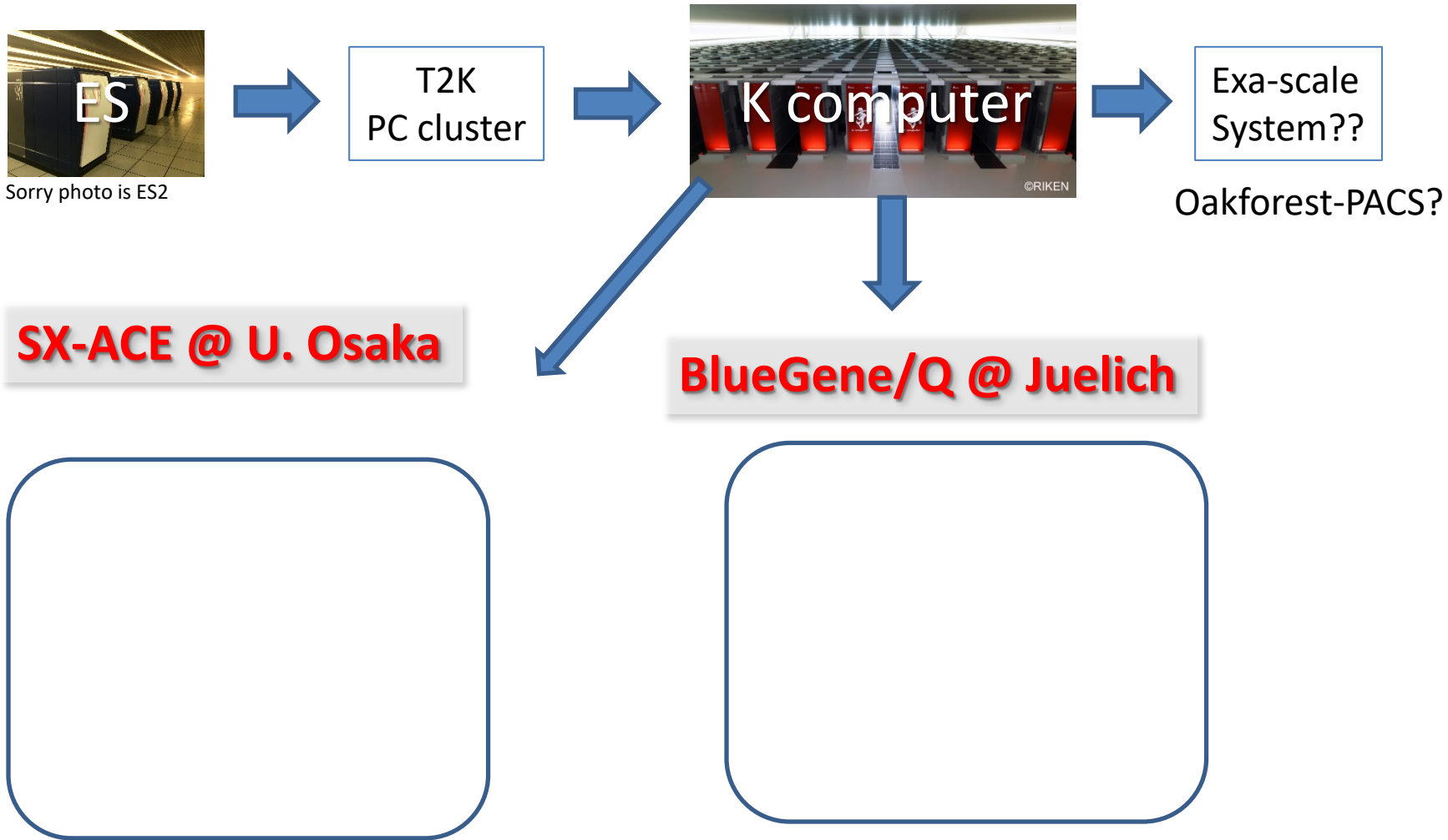
- **Development dedicated on K**
 - Almost hundred thousand proc. feasibility of algorithm and parallel implementation
 - Performance and scalability

Eigen-ES

2006

future of the EigenExa Project

- Porting the EigenExa library from K to other systems.



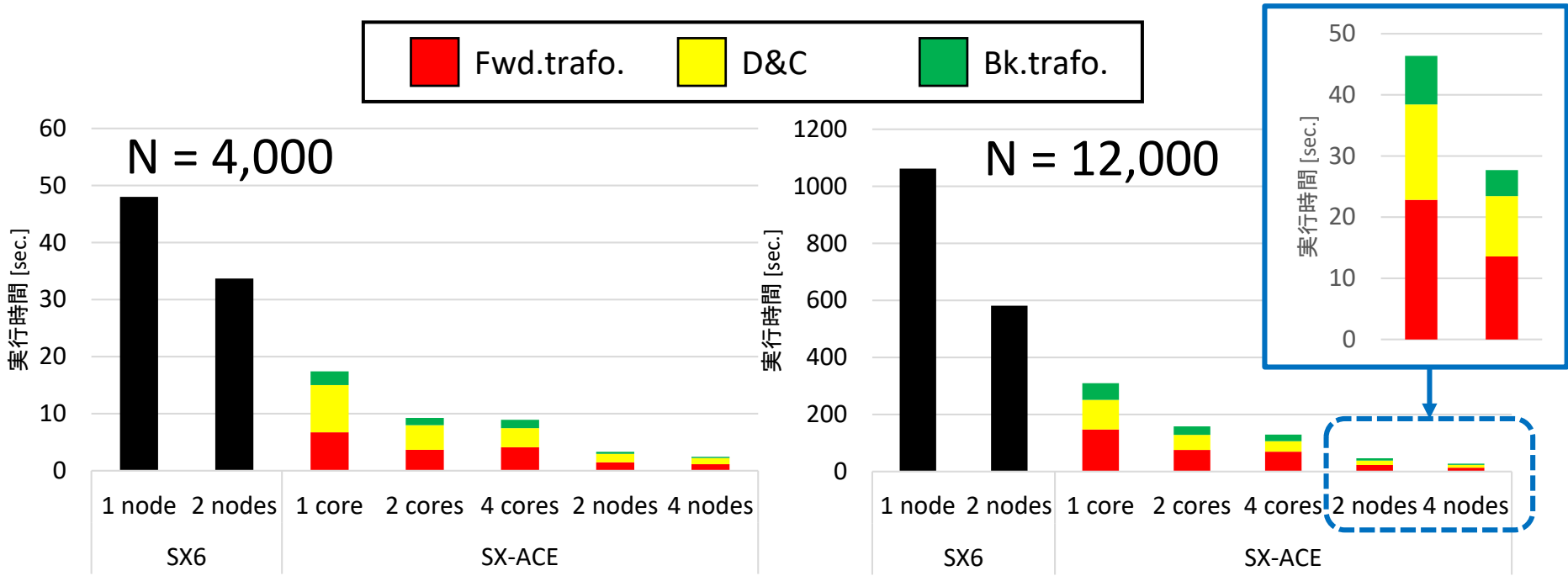
Hardware specification

	SX-ACE	SX-6 / ES1	京
FLOPS/node(PE or core)	64 GFLOPS/core	8 GFLOPS/PE	16 GFLOPS/core
Numbers of PEs	4 cores/node	8 PEs/node	8 cores/node
Memory bandwidth	256 GB/s	256 GB/s	64 GB/s
B/F	1	4	0.5
Network bandwidth (bi-direction)	8 GB/s	12.3 GB/s	10 GB/s

- SX-6/ES → SX-ACE
 - FLOPS/node increases 4 folds.
 - Relative bandwidth of memory and network degrade
- Next slide: Performance evaluation of eigensolver for a small problem
 - Randomly generated matrix
 - SX-ACE: EigenExa(eigen_sx), not optimized for SX-ACE
 - SX-6: the algorithm equivalent to eigen_s, 4PE/node was utilized.

Performance evaluation of eigensolver for a small problem

- SX-ACE: EigenExa(eigen_sx), not optimized for SX-ACE
- SX-6: the algorithm equivalent to eigen_s, 4PE/node was utilized.



- SX-6/SX-ACE = x12.5 when 2nodes and N=12,000
 - Hardware improvement + newly developed algorithm
- Peak performance when 4 nodes and N=12,000 reaches 32% of theoretical peak. (Fwd.trafo. 17%, D&C 33%, Bk.trafo. 79%)
- Even vector system, the forward transform is dominant.

Next Steps

- **Hardware**

- Near-future architecture, such as GPUs, MICs, FPGAs, accelerator boards, ...

- We always change and adapt the target architectures...

- for example, distributed parallel of multi-vector processors, on ES1

- the second generation was cluster of commodity processor and interconnect.

- present version is the third generation.

- **Target problems** (Complex, Tensor, Higher precision)

- Standard type eigenvalue problems is currently supposed.

- Generalized version is optional.

- Not only building IEEE754 double but wider format QP (quadruple precision) is being developed by taking advantage of double-double or multiple-double data format.

- **Algorithm** (revival of old but solid idea to post-Moore era's processing elements)

- Non-block algorithm but Tiling when we focus on local computing

- Hierarchical block strategy for a case of distributed computing

Target Architecture in near future

- We also have two branched projects from EigenExa on the K computer architecture
 - GPU:
 - Eigen-G = Experimental code on a single node + a single GPU environment
 - ASPEN.K2 = Automatic-tuning GPU BLAS kernels, especially, SYMV kernel
 - Intel Xeon Phi
 - Divide and conquer algorithm for GEVP focused on a pair of banded matrices
 - FPGAs ?



GPU

TI, etc, "Eigen-G: GPU-based eigenvalue solver for real-symmetric dense matrices", PPAM2013, LNCS8384
 TI, etc. "High Performance SYMV Kernel on a Fermi-core GPU", VECPAR 2012, LNCS 7851,
 TI, etc. "Automatic-tuning for CUDA-BLAS kernels by Multi-stage d-Spline Pruning Strategy", @²HPSC 2014

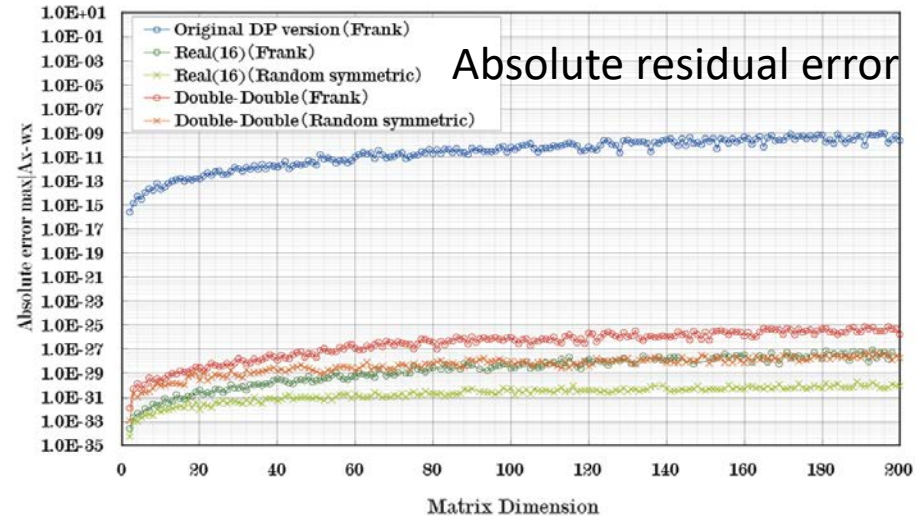


MIC

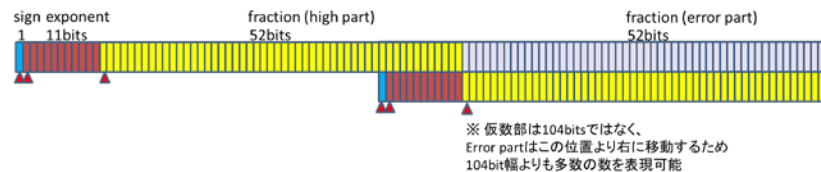
Y.Hirota, etc, "Divide-and-Conquer Method for Symmetric-Definite Generalized Eigenvalue Problems of Banded Matrices on Manycore Systems", SIAM LA15
 Y.Hirota, etc. "Acceleration of Divide and Conquer Method for Generalized Eigenvalue Problems of Banded Matrices on Manycore Architectures, PMAA14.

QP(Quadruple Precision)

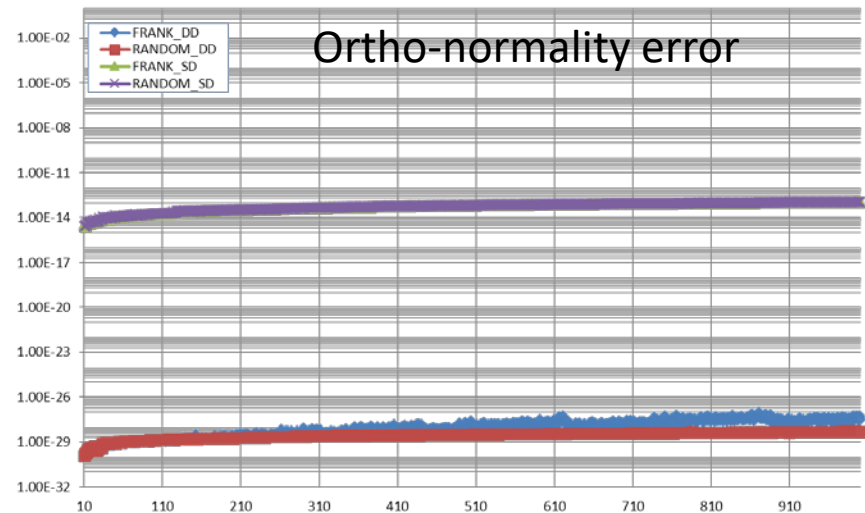
- Emerging long-time and large-scale computation, **rounding error** on the IEEE754 'double' floating point format with $O(10^{15})$ operations will be a considerable issue. The **DD, double-double, format** (D.H.Bailey, DDFUN90, <http://crd.ltdl.gov/~dhbailey/mpdist>) is one of promising technologies to ensure higher precision without the help of special hardware. The DD format consists of the 'high' and the 'error' parts, and their summation represents higher precision data.



- $$\mathbf{a}_{dd} := \mathbf{a}_{hi} + \mathbf{a}_{err}, \quad (|\mathbf{a}_{hi}| > |\mathbf{a}_{err}|)$$



- Addition and multiplication of two DD-format data are defined simply with approximately 20 double-precision floating operations. It is expected to help several issues on multicore platforms like accuracy and utilization problems. In this study, we are developing a double-double precision (**quadruple precision**) eigenvalue solver, 'QPEigenK'. It performs on distributed memory parallel computers. In addition, **OpenMP and MPI parallel models** are supported.



Y.Hirota, etc. HPC in Asia Award Winning Poster: Performance of Quadruple Precision Eigenvalue Solver Libraries QPEigenK & QPEigenG on the K Computer

Summary of talk

- **EigenExa project (2011-2016)**
 - The first milestone : 1million order eigenvalue computation with full nodes of K computer.
 - Second milestone : optimization of communication
- **We struggled against 3 walls of bottleneck**
 - **Memory bandwidth** → **Block algorithm**
 - **Network Latency** → **Communication avoiding (CA) and communication hiding (CH)**
 - **Parallelism** → **on-going work towards new hardware**
- **Near-Future work**
 - Establish the CA technology for total performance of EigenExa
 - Quadruple Precision version
 - Vector computers, other platforms
 - GPU cluster, MIC cluster, etc.
- **Topics for Collaboration is broad,**
 - **New target architectures, FPGA ? or ?**
 - **New topics must be also concerned like Reproducibility and FT**
 - **New Collaboration with CS and Applications!**

THANKS!
ご清聴
ありがとう
ございました

The results of the present study were obtained in part using the K computer at RIKEN Advanced Institute for Computational Science.

