



### The HPC Challenge (HPCC) Benchmark Suite

Characterizing a system with several specialized kernels

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"The HPC Challenge (HPCC) Benchmark Suite"

SC06, Tampa, Florida, Sunday, November 12, 2006

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- Opinions, interpretations, conclusions, and recommendations are those of the authors and are not necessarily endorsed by the United States Government

## Outline

- Overview
- The HPCC kernels
- Database & output formats
- HPCC award
- Augmenting TOP500
- Balance Analysis
- Conclusions

## Introduction

- Overview
  The Kernels
  Output formats
  HPCC awards
  Augm. TOP500
  Balance Analys
- Conclusions

- HPC Challenge Benchmark Suite
  - To examine the performance
     of HPC architectures using kernels
     with more *challenging* memory access patterns
     than HPL
  - To augment the TOP500 list
  - To provide benchmarks that *bound* the performance of many real applications as a function of memory access characteristics — e.g., spatial and temporal locality

## **TOP500 and HPCC**

- TOP500
  - Performance is represented by only a single metric
  - Data is available for an extended time period (1993-2006)
- Problem: There can only be one "*winner*"
- Additional metrics and statistics
  - Count (single) vendor systems on each list
  - Count total flops on each list per vendor
  - Use external metrics: price, ownership cost, power, ...
  - Focus on growth trends over time

- HPCC
  - Performance is represented by multiple single metrics
  - Benchmark is new so data is available for a limited time period (2003-2007)
- Problem: There cannot be one "*winner*"
- We avoid "*composite*" benchmarks
  - Perform trend analysis
    - HPCC can be used to show complicated kernel/ architecture performance characterizations
  - Select some numbers for comparison
  - Use of kiviat charts
    - Best when showing the differences due to a single independent "variable"
  - Compute balance ratios
- Over time also focus on growth trends

# High Productivity Computing Systems (HPCS)

#### Goal:

Provide a new generation of economically viable high productivity computing systems for the national security and industrial user community (2010)

#### Impact:

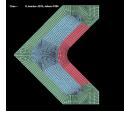
- **Performance** (time-to-solution): speedup critical national security applications by a factor of 10X to 40X
- **Programmability** (idea-to-first-solution): reduce cost and time of developing application solutions
- **Portability** (transparency): insulate research and operational application software from system
- Robustness (reliability): apply all known techniques to protect against outside attacks, hardware faults, & programming errors

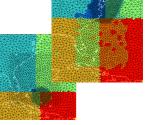


#### **Applications:**









Analysis & Assessment

Software Technology

Analysis & Assessmen

**HPCS Program Focus Areas** 

Industry R&D

R&D

ogramming

Hardware

Technology

Industry

Performance

Characterizatior & Prediction

System

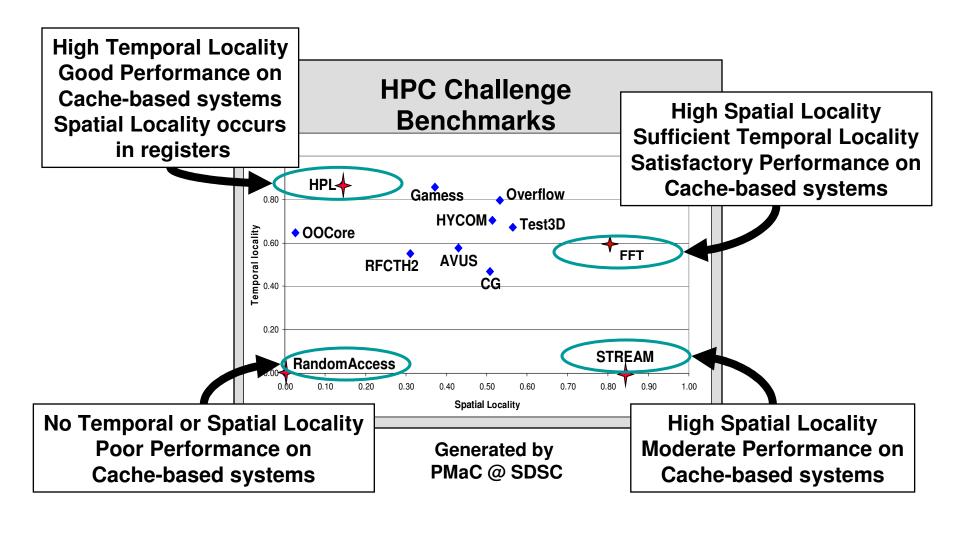
Architecture



 Intelligence/surveillance, reconnaissance, cryptanalysis, weapons analysis, airborne contaminant modeling and biotechnology

Fill the Critical Technology and Capability Gap Today (late 80's HPC technology).....to.....Future (Quantum/Bio Computing)

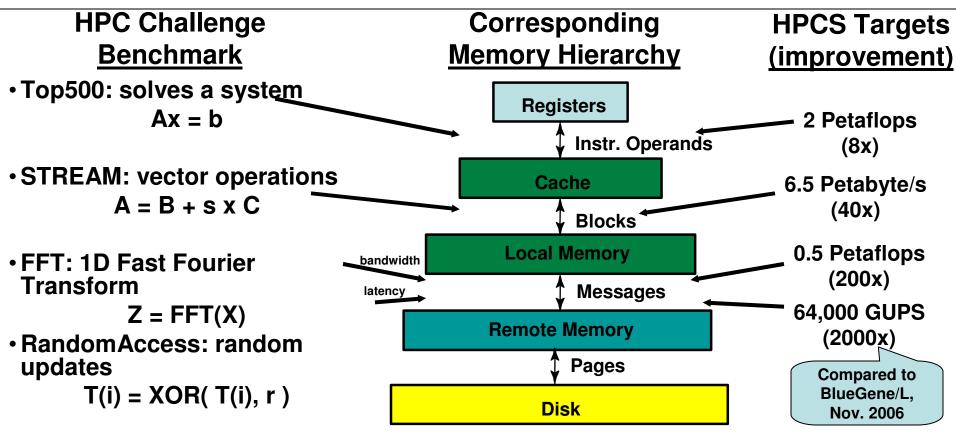
### Motivation of the HPCC Design



Spatial and temporal data locality here is for one node/processor — i.e., locally or "in the small"

Further information: "Performance Modeling and Characterization" @ San Diego Supercomputer Center http://www.sdsc.edu/PMaC/ 7/31

### **HPCS Performance Targets**



- HPCS program has developed a new suite of benchmarks (HPC Challenge)
- Each benchmark focuses on a different part of the memory hierarchy
- HPCS program performance targets will flatten the memory hierarchy, improve real application performance, and make programming easier

## HPCC as a Framework (1/2)

- Many of the component benchmarks were widely used before
  - HPCC is more than a packaging effort
  - E.g., provides consistent verification and reporting
- Important:

Running these benchmarks on a single machine — with a single configuration and options

- The benchmark components are still useful separately for the HPC community, meanwhile
- The unified HPC Challenge framework creates an unprecedented view of performance characterization of a system
  - A comprehensive view with data captured under the same conditions allows for a variety of analyses depending on end user needs

## HPCC as a Framework (2/2)

- A single executable is built to run all of the components
  - Easy interaction with batch queues
  - All codes are run under the same OS conditions just as an application would
    - No special mode (page size, etc.) for just one test (say Linpack benchmark)
    - Each test may still have its own set of compiler flags
      - Changing compiler flags in the same executable may inhibit interprocedural optimization
- Scalable framework Unified Benchmark Framework
  - By design, the HPC Challenge Benchmarks are scalable with the size of data sets being a function of the largest HPL matrix for the tested system

## **HPCC Tests at a Glance**

#### • Overview

•

- The Kernels
- Output formats
- HPCC awards
- Augm. TOP500
- Balance Analys.
- Conclusions

High Performance LinpackSolving Ax = b $A \in \mathbb{R}^{n \times n}$  $x, b \in \mathbb{R}$ 

### 2. DGEMM

1. HPL

- Double-precision General Matrix-matrix Multiply
- **Computing**  $C \leftarrow \alpha AB + \beta C$   $A, B, C \in \mathbb{R}^{n \times n}$   $\alpha, \beta \in \mathbb{R}$
- <u>Temporal/spatial locality:</u> similar to HPL

### 3. STREAM

- measures sustainable memory bandwidth with vector operations
- COPY: c=a SCALE:  $b=\alpha c$ ADD: c=a+b TRIAD:  $a=b+\alpha c$

### 4. PTRANS

- Parallel matrix TRANSpose
- **Computing**  $A = A^T + B$
- <u>Temporal/spatial locality:</u> similar to EP-STREAM, but includes global communication

### 5. RandomAccess

 calculates a series of integer updates to random locations in memory

```
Ran = 1;
for (i=0; i<4*N; ++i) {
    Ran= (Ran<<1) ^
        (((int64_t)Ran < 0) ? 7:0);
    Table[Ran & (N-1)] ^= Ran;
}
```

- Use at least 64-bit integers
- About half of memory used for 'Table'
- Parallel look-ahead limited to 1024

### 6. FFT

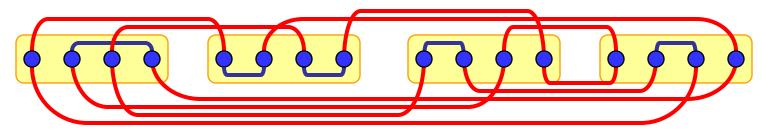
- Fast Fourier Transform
- **Computing**  $z_k = \Sigma x_j \exp(-2\pi \sqrt{-1} jk/n)$   $x, z \in \mathbb{C}^n$

### 7. b\_eff

- Patterns: ping-pong,
  - natural ring, and
  - random ring patterns
- Bandwidth (w 2,000,000 bytes messages)
- Latency (with 8 bytes messages)

## **Random Ring Bandwidth**

- Reflects communication patterns in unstructured grids
- And 2<sup>nd</sup> & 3<sup>rd</sup> dimension of a Cartesian domain decomposition
- On clusters of SMP nodes:
  - Some connections are inside of the nodes
  - <u>Most</u> connections are inter-node

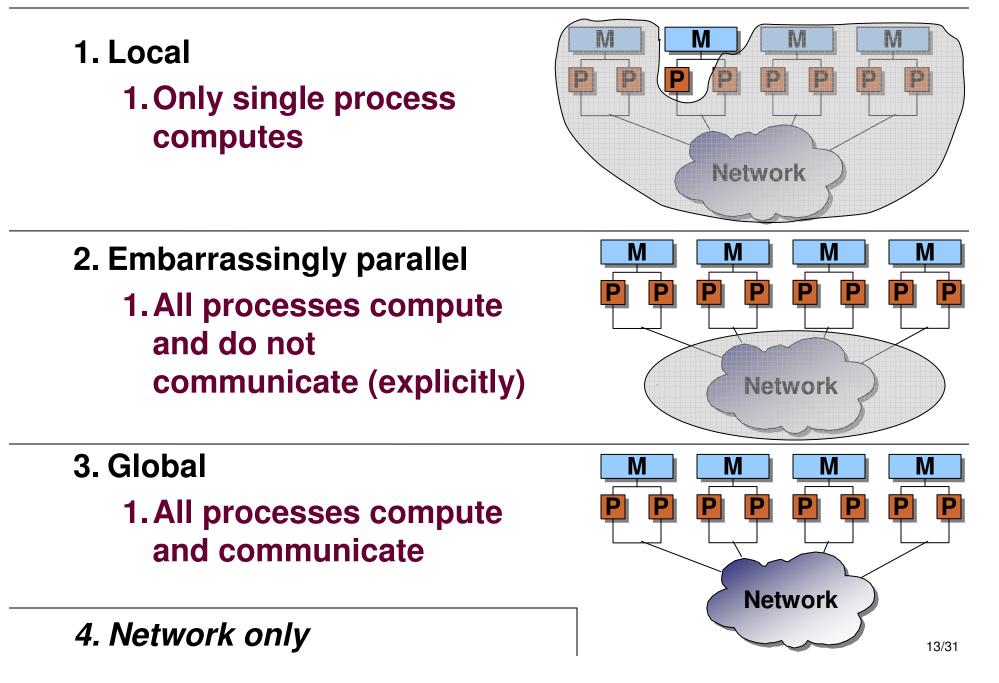


- Global benchmark  $\rightarrow$  all processes participate
- Reported: bandwidth per process
- Accumulated bandwidth

:= bandwidth per process x #processes



### **HPCC Testing Scenarios**



### **Base vs. Optimized Submission**

G-RandomAccess

			$\frown$			$\frown$						
System Ir	nformation		Run	G-HPL	G-PTRANS	G-Randon Access	G-FFTE	G-STREAM Triad	EP STREAM Triad	EP DGEMM	Random Ring Bandwidth	Random Ring Latency
System - Processor	Speed	Count	Туре	TFlop/s	GB/s	Gup/s	GFlop/s	GB/s	GB/s	GFlop/s	GB/s	usec
Cray mfeg8 X1E	1.13GHz	248	opt	3.3889	66.01	1.85475	-1	3280.9	13.229	13.564	0.29886	14.58
Cray X1E X1E MSP	1.13GHz	252	base	3.1941	85.204	0.014868	15.54	2440	9.682	14.185	0.36024	14.93

- Base code: <u>Latency</u> based execution
- Optimization I: <u>UPC based code only a few lines</u>
  - Optimization inside of UPC compiler / library
  - ~125x improvement
- Optimization II: <u>Butterfly (MPI-based) algorithm</u>
  - Bandwidth based (packet size ~ 4 kB)
  - On BlueGene/L with special communication library: 537 x faster than "base"

### **Results**

Overview

- The Kernels
- Output formats
- HPCC awards
- Augm. TOP500
- Balance Analys.
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#### HPCC Database

### ←upload of HPCC results

### $\rightarrow$ Output through several interfaces

- Web-output
  - Table with several subsets of kernels

#### - Base / optimized / base+optimized

#### Can be sorted by any column

Condensed Results - Base Runs Only - 132 Systems - Generated on Tue Jun 19 08:54:47 2007

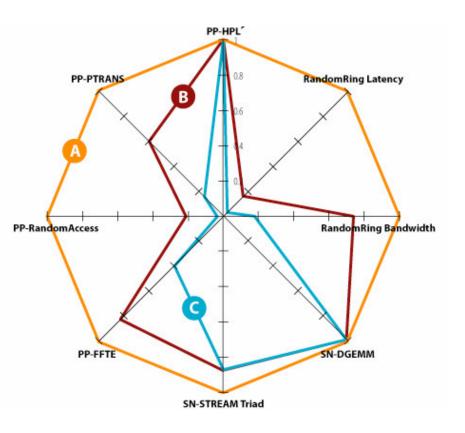
System Informa System - Processor - Speed - Coun		ses		G-HPL	G-PTRANS	G-Random Access	G-FFTE	EP-STREAM Sys	EP-STREA Triad	EP-DGEMM	Random Ring Bandwidth	RandomRing Latency
MA/PT/PS/PC/TH/PR/CM	CS/IC/IA/SD			TFlop/s	GB/s	Gup/s	GFlop/s	GB/s	GB/s	GFlop/s	GB/s	usec
Cray Inc. Red Storm/XT3 AMD Opteron	2.4GHz	12960	125920	91.0350000	2356.9700	1.7401500	1554.0700	54840.499	2.1158	4.39939	0.05911	16.29
IBM Blue Gene/L PowerPC 440	0.7GHz	65536	165536	80.6830000	339.2840	0.0657312	2178.1100	53555.888	0.8172	1.85619	0.01084	8.84
IBM p5-575 Power5	1.9GHz	10240	110240	57.8670000	553.0090	0.1693440	842.5000	55184.179	5.3891	7.08562	0.11015	118.59
IBM p5-575 Power5	1.9GHz	8192	1 8192	45.7019000	2626.1700	0.3239760	908.6920	44455.936	5.4268	7.06423	0.08871	11.05
Cray Inc. XT3 Dual-Core AMD Opteron	2.6GHz	10404	110404	43.4033000	778.3850	0.8235630	1107.2100	25774.557	2.4774	4.78995	0.06937	14.32
IBM Blue Gene/L PowerPC 440	0.7GHz	65536	165536	37.3540000	4665.9100	0.1648600	1762.8200	62889.787	0.9596	2.47017	0.01039	8.62
Cray Inc. XT3 AMD Opteron	2.6GHz	8190	1 8190	35.1985000	603.1050	0.7308180	882.4230	17998.835	2.1977	4.79150	0.08599	14.22
	1.000	04.00		00.0175000	575 0000	0.0000000	000 0000	40000 400	5 0 4 7 0	0.00010	0.07000	51.00

- As Excel or XML (all results)
- Comparing up to 6 platforms with a Kiviat diagram

## **Kiviat Charts: Comparing Interconnects**

- Comparing per-process values
- 8 fixed benchmark kernels
- Up to 6 systems
- Normalized:
  - 1 = best system at each kernel
- Example:
  - AMD Opteron clusters
    - 2.2 GHz
    - 64-processor cluster
  - Interconnects
    - 1. GigE
    - 2. Commodity
    - 3. Vendor
  - Cannot be differentiated based on:
    - HPL
    - Matrix-matrix multiply
- Available on HPCC website

#### Kiviat chart (radar plot)



## **HPCC Awards Overview**

- Goals
  - Increase awareness of HPCC benchmarks
  - Increase awareness of HPCS program and its goals
  - Increase number of HPCC submissions
    - Expanded view of largest supercomputing installations
- Means
  - HPCwire sponsorships and press coverage
  - HPCS mission partners' contribution
  - HPCS vendors' contribution
- Awards are presented at the SCxx HPC Challenge BOF

• The Kernels

Overview

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### **HPCC Awards Rules**

- Class 1: Best Performance
  - Figure of merit: raw system performance
  - Submission must be valid HPCC database entry
    - Side effect: populate HPCC database
  - 4 categories: HPCC components
    - HPL
    - STREAM-system
    - RandomAccess
    - FFT
  - Award certificates
    - 4x \$500 from HPCwire

### Class 2: Most Productivity

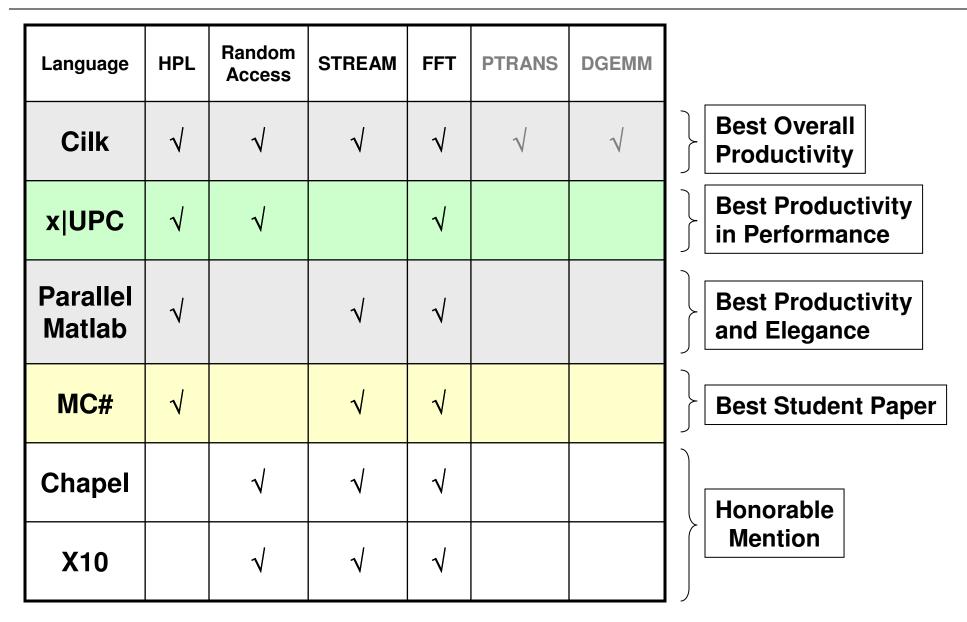
- Figure of merit: performance (50%) and elegance (50%)
  - Highly subjective
  - Based on committee vote
- Submission must implement at least 3 out of 4 Class 1 tests
  - The more tests the better
- Performance numbers are a plus
- The submission process:
  - Source code
  - "Marketing brochure"
  - SC06 BOF presentation
- Award certificate
  - \$1500 from HPCwire

## HPC #

### SC|06 HPCC Award – Class 1

<u>G-HPL</u>	Achieved	System	Affiliation	Submitter
1st place	259 Tflop/s	IBM BG/L	DOE/NNSA/LLNL	Tom Spelce
1st runner up	67 Tflop/s	IBM BG/L	IBM T.J. Watson	John Gunnels
2nd runner up	57 Tflop/s	IBM p5-575	LLNL	Charles Grassl
HPCS goal:	2000 Tflop/s	= current 1st p	lace <u>x 8</u>	
<b>EP-STREAM-Triad</b>	Achieved	System	Affiliation	Submitter
1st place	160 TB/s	IBM BG/L	DOE/NNSA/LLNL	Tom Spelce
1st runner up	55 TB/s	IBM p5-575	LLNL	Charles Grassl
2nd runner up	43 TB/s	Cray XT3	SNL	Courtenay Vaughan
HPCS goal:	6500 TB/s	= current 1st p	lace <u>x 40</u>	
<u>G-FFT</u>	Achieved	System	Affiliation	Submitter
<u>G-FFT</u> 1st place	Achieved 2.311 Tflop/s	System IBM BG/L	Affiliation DOE/NNSA/LLNL	Submitter Tom Spelce
1st place	2.311 Tflop/s	IBM BG/L	DOE/NNSA/LLNL	Tom Spelce
1st place 1st runner up	<b>2.311 Tflop/s</b> 1.122 Tflop/s	IBM BG/L Cray XT3 Dual Cray XT3	DOE/NNSA/LLNL ORNL SNL	Tom Spelce Jeff Larkin
1st place 1st runner up 2nd runner up	2.311 Tflop/s 1.122 Tflop/s 1.118 Tflop/s 500.0 Tflop/s	IBM BG/L Cray XT3 Dual Cray XT3	DOE/NNSA/LLNL ORNL SNL	Tom Spelce Jeff Larkin
1st place 1st runner up 2nd runner up <b>HPCS goal:</b>	2.311 Tflop/s 1.122 Tflop/s 1.118 Tflop/s 500.0 Tflop/s	IBM BG/L Cray XT3 Dual Cray XT3 <b>= current 1st p</b>	DOE/NNSA/LLNL ORNL SNL Iace <u>x 200</u>	Tom Spelce Jeff Larkin Courtenay Vaughan
1st place 1st runner up 2nd runner up <b>HPCS goal:</b> <b>G-RandomAccess</b>	2.311 Tflop/s 1.122 Tflop/s 1.118 Tflop/s 500.0 Tflop/s Achieved	IBM BG/L Cray XT3 Dual Cray XT3 <b>= current 1st p</b> System	DOE/NNSA/LLNL ORNL SNL lace <u>x 200</u> Affiliation	Tom Spelce Jeff Larkin Courtenay Vaughan <b>Submitter</b>
1st place 1st runner up 2nd runner up <b>HPCS goal:</b> <u><b>G-RandomAccess</b></u> 1st place	2.311 Tflop/s 1.122 Tflop/s 1.118 Tflop/s 500.0 Tflop/s Achieved 35 GUPS	IBM BG/L Cray XT3 Dual Cray XT3 <b>= current 1st p</b> System IBM BG/L	DOE/NNSA/LLNL ORNL SNL lace <u>x 200</u> Affiliation DOE/NNSA/LLNL	Tom Spelce Jeff Larkin Courtenay Vaughan <b>Submitter</b> Tom Spelce

### SC|06 HPCC Awards Class 2



# Augmenting TOP500's 26<sup>th</sup> Edition

<u>Augm. TOP500</u>
Balance Analys.
Conclusions

. . .

								• Conclu	510115
	Computer	Rmax	HPL	PTRANS	STREAM	FFT	GUPS	Latency	B/W
1	BlueGene/L	281	259	374	160	2311	35.5	6	0.2
2	BGW (**)	91	84	172	50	1235	21.6	5	0.2
3	ASC Purple	63	58	576	44	967	0.2	5	3.2
4	Columbia (**)	52	47	91	21	230	0.2	4	1.4
5	Thunderbird	38							
6	Red Storm	36	33	1813	44	1118	1.0	8	1.2
7	Earth Simulator	36							
8	MareNostrum	28							
9	Stella	27							
10	Jaguar	21	20	944	29	855	0.7	7	<b>1.2</b>

# Augmenting TOP500's 28<sup>th</sup> Edition with HPCC

1

2

3

4

5

6

G-EP-G-G-G-PingPong PingPong Rmax Random Computer HPL PTRANS STREAM FFT Latency Bandwidth Access TFlop/s GB/s μs TFlop/s GB/s Triad TB/s GFlop/s GUPS 259.2 4665.9 160 35.47 2311 5.92 µs 0.158 BlueGene/L 280.6 80.7 339.3 57 0.066 0.157 2178 **7.07** μs 29.81 Cray XT3 Red Storm Opteron 101.4 91.0 2357.0 55 1554 7.16 us 2.024 dual-core 1.74 83.9 \*\* BlueGene/L 171.55 \*\* 50 \*\* **1235**\*\* **21.61**\*\* BGW IBM/Watson 91 0.159 4.95 us 39 \*\* 1391\*\* 0.348\*\* 109 \*\* 37 \*\* (\*\* 32768→ 40960) ASC Purple <sup>IBM</sup> <sub>p5</sub> 1.02(\*) 69 \*\* 66 \*\* 1004\*\* 659 \*\* 75.8 5.10 µs 3.184 0.202\*\* (\*\* 10240→ 12208) Upper values 62.63 = optimized MareNostrum Bottom values = base Thunderbird 53.00 (\*) = not published in HPCC database FO 04

7	lera-10	52.84			<u>= extrapol</u>	ated: #C	PUs HPC	C → #CPl	Js Linpack
8	Columbia <sup>SGI Altix</sup> Infiniband (** 2024→ 10160)	51.87	47 **	91.31 **	20 **	229 **	0.25 **	4.23 μs	0.896
9	TSUBAME	47.38							
10	Cray XT3 Jaguar Opteron dual-core	43.48	43.40	<mark>2039</mark> 778	27	1127 1107	10.67 0.82	6.69 µs	<b>1.15</b>

# Augmenting TOP500's 28<sup>th</sup> Edition with HPCC

1

2

3

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10

**TSUBAME** 

Jaguar

Crav XT3

Opteron dual-core 47.38

43.48

43.40

Random Ring BW Random Pina Ping HPL **STREAM** Rmax Computer global per proc. Pona Rina Pona TFlop/s TFlop/s **Triad TB/s** TB/s GB/s GB/s Latency µs 259.2 0.727 160 0.011 0.158 7.78 5.92 BlueGene/L 280.6 80.7 0.710 0.011 57 0.157 8.84 7.07 Cray XT3 Red Storm Opteron dual-core 101.4 91.0 1.532 2.024 55 0.059 16.29 7.16 BlueGene/L 83.9 BGW IBM/Watson 171.55 0.490 \*\* 91 0.012 0.159 9.51 4.95 39 \*\* 109 \*\* (\*\* 32768→ 40960) ASC Purple IBM p5 HPS 1.345 \*\* 69 \*\* 66 \*\* 75.8 0.110 3.154 118.59 5.10 (\*\* 10240→ 12208) MareNostrum 62.63 B/W Latency **Global values** Random / PingPong ratio Thunderbird 53.00 (i.e., accumulated ratio: 1.3 – 2.3: PingPing / Random per system) 23 (Purple) 7 - 34 Tera-10 52.84 Columbia SGI Altix 47 \*\* 20 \*\* 1.247 \*\* 51.87 0.122 0.896 6.98 4.23 (\*\* 2024→ 10160)

Upper values

**Bottom values** 

27

= optimized

0.069

= base

0.722

23/31

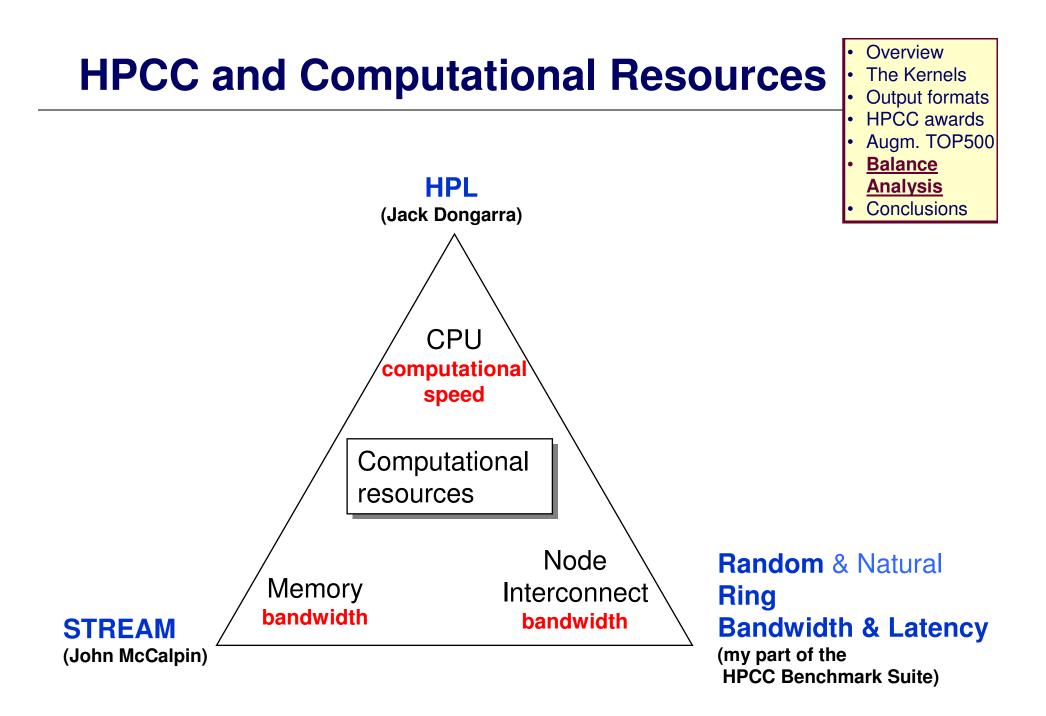
6.69

(\*) = not published in HPCC DB

14.32

(\*\*) = extrapolated

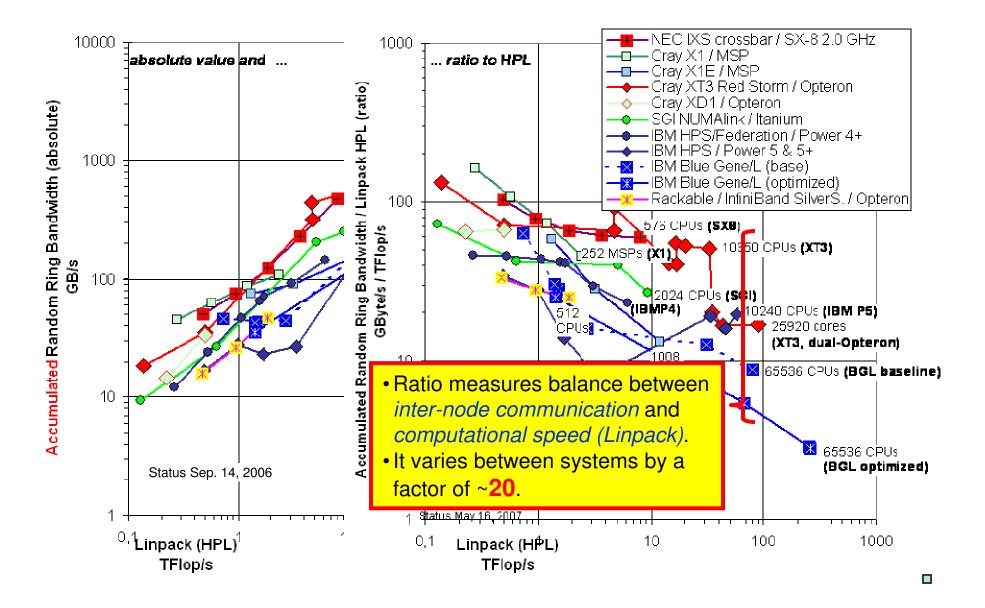
1.15



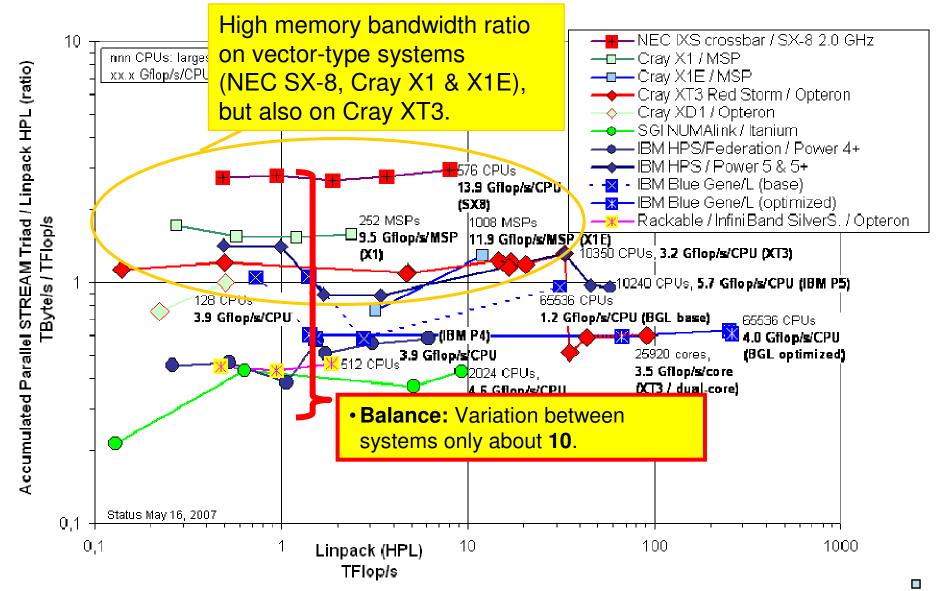
### **Balance Analysis with HPCC Data**

- Balance can be expressed as a set of ratios •
  - e.g., accumulated memory bandwidth / accumulated **Tflop/s rate**
- Basis
  - Linpack (HPL) → Computational Speed
  - − Random Ring Bandwidth → Inter-node communication
- - Parallel STREAM Copy or Triad → Memory bandwidth
- Be careful: •
  - Balance calculation always with accumulated data on the total system (Global or EP)
  - Random Ring B/W: per process value must be multiplied by *#processes*

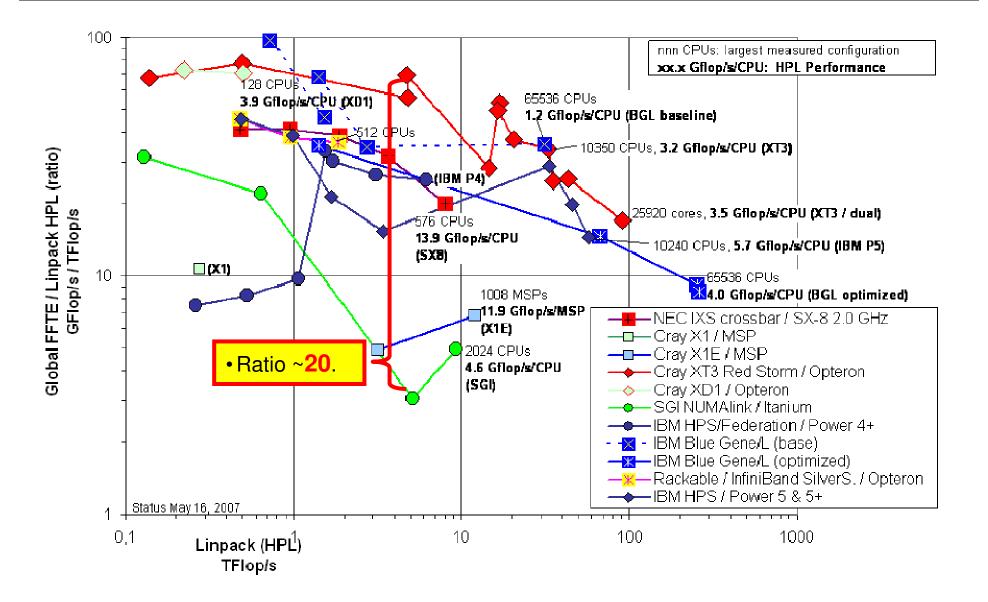
### **Balance: Random Ring B/W and HPL**



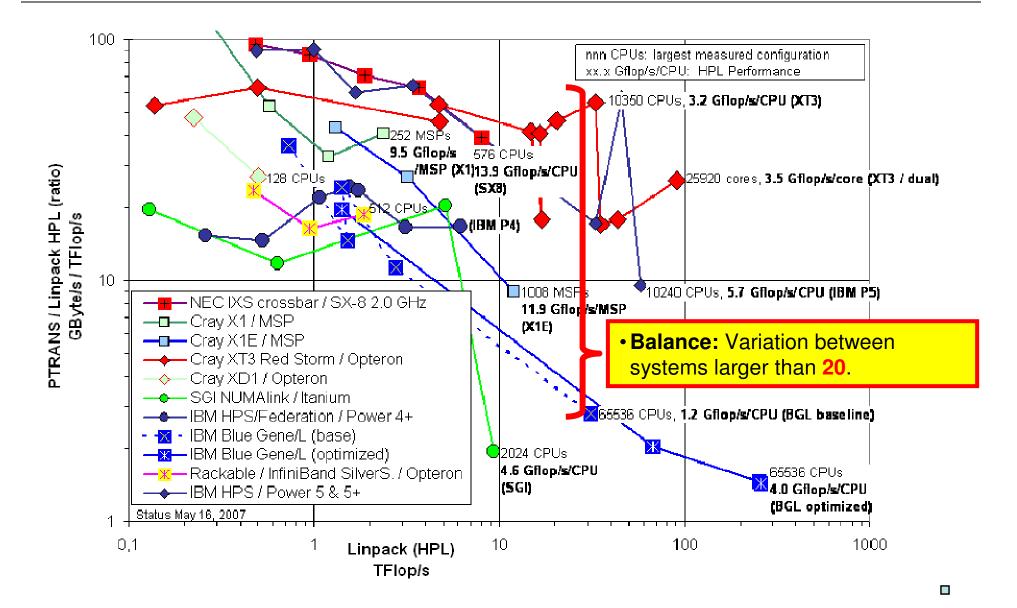
### **Balance: Memory and CPU Speed**



### **Balance: FFT and CPU**



### **Balance: PTRANS and CPU**



### **Acknowledgments**

- Thanks to
  - all persons and institutions that have uploaded HPCC results.
  - Jack Dongarra and Piotr Luszczek for inviting me into the HPCC development team.
  - Matthias Müller, Sunil Tiyyagura and Holger Berger for benchmarking on the SX-8 and SX-6 and discussions on HPCC.
  - Nathan Wichmann from Cray for additional Cray XT3 and X1E data.

### References

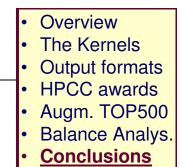
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   Proceedings of the 12th European PVM/MPI Users' Group Meeting, EuroPVM/MPI 2005
- <u>http://icl.cs.utk.edu/hpcc/</u>

## Conclusions

- HPCC is an interesting basis for
  - benchmarking computational resources
  - Augmenting TOP500
  - analyzing the balance of a system
  - scaling with the number of processors
  - with respect to applications' needs (e.g., locality characteristics)
- HPCC helps to show the strength and weakness of supercomputers
- Future super computing should not focus only on Pflop/s in the TOP500
  - Memory and network bandwidth are as same as important to predict real application performance

Copy of the slides:

http://www.hlrs.de/people/rabenseifner/publ/publications.html#SPEC2007



### Appendix

### **HPCC Tests - HPL**

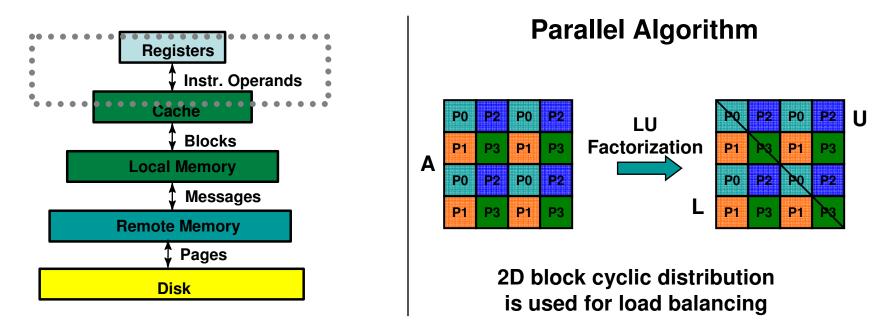
- HPL = High Performance Linpack
- <u>Objective</u>: solve system of linear equations

 $Ax=b \qquad A \in \mathbf{R}^{n \times n} \qquad x, b \in \mathbf{R}$ 

- <u>Method</u>: LU factorization with partial row pivoting
- **Performance**:  $(\frac{2}{3}n^3 + \frac{3}{2}n^2)/t$
- <u>Verification</u>: scaled residuals must be small  $|| Ax-b || / (\varepsilon ||A|| ||x|| n)$
- <u>Restrictions</u>:
  - No complexity reducing matrix-multiply
    - (Strassen, Winograd, etc.)
  - 64-bit precision arithmetic through-out
    - (no mixed precision with iterative refinement)

## **HPCC HPL: Further Details**

- High Performance Linpack (HPL) solves a system Ax = b
- Core operation is a LU factorization of a large MxM matrix
- Results are reported in floating point operations per second (flop/s)



- Linear system solver (requires all-to-all communication)
- Stresses local matrix multiply performance
- DARPA HPCS goal: 2 Pflop/s (8x over current best)

## HPCC Tests - DGEMM

- DGEMM = Double-precision General Matrix-matrix Multiply
- Objective: compute matrix
  - $C \leftarrow \alpha AB + \beta C \qquad A, B, C \in \mathbb{R}^{n \times n} \quad \alpha, \beta \in \mathbb{R}$
- <u>Method</u>: standard multiply (maybe optimized)
- **Performance**: 2n<sup>3</sup>/t
- <u>Verification</u>: Scaled residual has to be small  $||x-y||/(\varepsilon n ||y||)$

where x and y are vectors resulting from multiplication by a random vector of left and right hand size of the objective expression

- <u>Restrictions</u>:
  - No complexity reducing matrix-multiply
    - (Strassen, Winograd, etc.)
  - Use only 64-bit precision arithmetic
- <u>Temporal/spatial Locality</u>: similar to HPL

## **HPCC Tests - STREAM**

- STREAM is a test that measures sustainable memory bandwidth (in Gbyte/s) and the corresponding computation rate for four simple vector kernels
- **Objective:** set a vector to a combination of other vectors

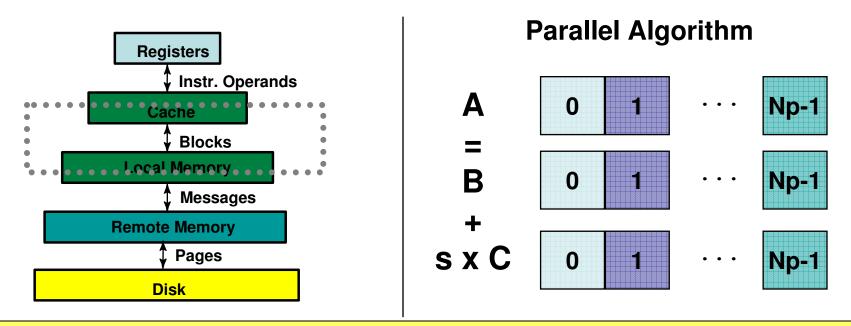
COPY:	c = a
SCALE:	$b = \alpha c$
ADD:	c = a + b
TRIAD:	$a = b + \alpha c$

- <u>Method</u>: simple loop that preserves the above order of operations
- **Performance**: 2n/t or 3n/t
- <u>Verification</u>: scalre residual of computed and reference vector needs to be small
   || x - y || / (ε n || y || )
- <u>Restrictions</u>:

- Use only 64-bit precision arithmetic

### **HPCC STREAM: Further Details**

- · Performs scalar multiply and add
- Results are reported in bytes/second



- Basic operations on large vectors (requires no communication)
- Stresses local processor to memory bandwidth
- DARPA HPCS goal: 6.5 Pbyte/s (40x over current best)

# **HPCC Tests - PTRANS**

- PTRANS = Parallel TRANSpose
- <u>Objective</u>: update matrix with sum of its transpose and another matrix

 $A = A^T + B \qquad A, B \in \mathbb{R}^{n \times n}$ 

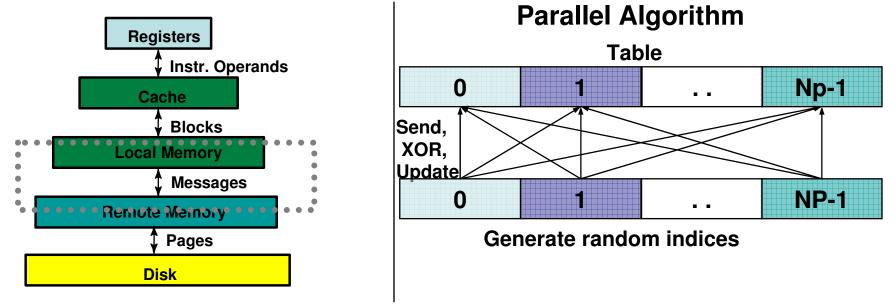
- <u>Method</u>: standard distributed memory algorithm
- **Performance**:  $n^2/t$
- Verification: scaled residual between computed and reference matrix needs to be small
   || A<sub>0</sub> – A || / (ε n || A<sub>0</sub> || )
- <u>Restrictions</u>:
  - Use only 64-bit precision arithmetic
  - The same data distribution method as HPL
- <u>Temporal/spatial Locality</u>: similar to EP-STREAM, but includes global communication

# **HPCC Tests - RandomAccess**

- RandomAccess calculates a series of integer updates to random locations in memory
- Objective: perform computation on Table
  Ran = 1;
  for (i=0; i<4\*N; ++i) {
   Ran= (Ran<<1) ^ (((int64\_t)Ran < 0) ? 7:0);
   Table[Ran & (N-1)] ^= Ran;
  }</pre>
- <u>Method</u>: loop iterations may be independent
- **Performance**: 4N/t
- <u>Verification</u>: up to 1% of updates can be incorrect
- <u>Restrictions</u>:
  - Use at least 64-bit integers
  - About half of memory used for 'Table'
  - Parallel look-ahead limited to 1024 (limit locality)

# **HPCC RandomAccess: Further Details**

- Randomly updates N element table of unsigned integers
- Each processor generates indices, sends to all other processors, performs XOR
- Results are reported in Giga Updates Per Second (GUPS)



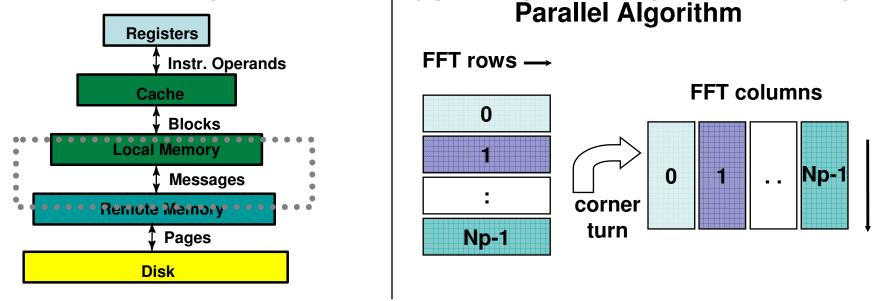
- Randomly updates memory (requires all-to-all communication)
- Stresses interprocessor communication of *small* messages
- DARPA HPCS goal: 64,000 GUPS (2000x over current best)

# **HPCC Tests - FFT**

- FFT = Fast Fourier Transform
- <u>Objective</u>: compute discrete Fourier Transform  $z_k = \sum x_j \exp(-2\pi \sqrt{-1} jk/n)$   $x, z \in \mathbb{C}^n$
- <u>Method</u>: any standard framework (maybe optimized)
- <u>Performance</u>: 5nlog<sub>2</sub>n/t
- <u>Verification</u>: scaled residual for inverse transform of computed vector needs to be small  $||x x^{(0)}|| / (\epsilon \log_2 n)$
- <u>Restrictions</u>:
  - Use only 64-bit precision arithmetic
  - Result needs to be in-order (not bit-reversed)

# **HPCC FFT: Further Details**

- 1D Fast Fourier Transforms an N element complex vector
- Typically done as a parallel 2D FFT
- Results are reported in floating point operations per second (flop/s)



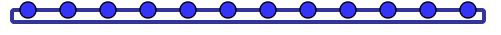
- FFT a large complex vector (requires all-to-all communication)
- Stresses interprocessor communication of *large* messages
- DARPA HPCS goal: 0.5 Pflop/s (200x over current best)

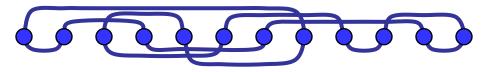
# HPCC Tests – b\_eff

- b\_eff measures effective bandwidth and latency of the interconnect
- <u>Objective</u>: exchange 8 (for latency) and 2000000 (for bandwidth) messages in
  - ping-pong,
  - natural ring, and
  - random ring patterns
- <u>Method</u>: use standard MPI point-to-point routines
- <u>Performance</u>: n/t (for bandwidth)
- <u>Verification</u>: simple checksum on received bits
- <u>Restrictions</u>:
  - The messaging routines have to conform to the MPI standard

# **HPCC b\_eff: Further Details**

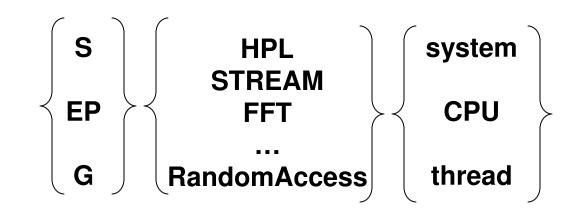
- Parallel communication pattern on <u>all</u> MPI processes:
  - Natural ring
  - Random ring

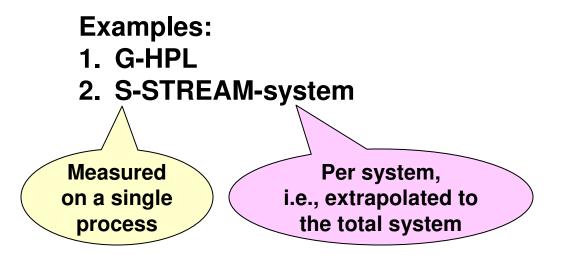




- Bandwidth per process
  - Accumulated message size / wall-clock time / number of processes
  - On each connection messages in both directions
  - − With 2xMPI\_Sendrecv and MPI non-blocking  $\rightarrow$  best result is used
  - Message size = 2,000,000 bytes
- Latency
  - Same patterns, message size = 8 bytes
  - Wall-clock time / (number of sendrecv per process)

## **Naming Conventions**





## **Base vs. Optimized Run**

- HPC Challenge encourages users to develop optimized benchmark codes that use architecture specific optimizations to demonstrate the best system performance
- Meanwhile, we are interested in both
  - The base run with the provided reference implementation
  - An optimized run
- The base run represents behavior of legacy code because
  - It is conservatively written using only widely available programming languages and libraries
  - It reflects a commonly used approach to parallel processing sometimes referred to as hierarchical parallelism that combines
    - Message Passing Interface (MPI)
    - OpenMP Threading
  - We recognize the limitations of the base run and hence we encourage optimized runs
- Optimizations may include alternative implementations in different programming languages using parallel environments available specifically on the tested system
- We require that the information about the changes made to the original code be submitted together with the benchmark results
  - We understand that full disclosure of optimization techniques may sometimes be impossible
  - We request at a minimum some guidance for the users that would like to use similar optimizations in their applications

## SC|05 HPCC Awards Class 2

Language	HPL	RandomAccess	STREAM	FFT	Sample submission from	
Python+MPI		$\checkmark$	√			
pMatlab	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$	<b>committee members</b>	
Cray MTA C		$\checkmark$		$\checkmark$		
UPCx3	$\checkmark$	$\checkmark$	$\checkmark$		Winners	
Cilk	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$	Finalists	
Parallel Matlab	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$		
MPT C	$\checkmark$			$\checkmark$		
OpenMP, C++		$\checkmark$	$\checkmark$			
StarP	$\checkmark$		$\checkmark$			
HPF	$\checkmark$			$\checkmark$		

#### Augmenting TOP500's 27<sup>th</sup> Edition with HPCC

June 2006 -

	Computer	Rmax	HPL	PTRANS	STREAM	FFT	GUPS	Latency	B/W
1	BlueGene/L	280.6	259.2	4665.9	160	2311	35.47	5.92	0.159
2	BGW (**)	91	83.9	171.55	50	1235	21.61	4.70	0.159
3	ASC Purple	75.8	57.9	553	55	842	1.03	5.1	3.184
4	Columbia (**)	51.87	46.78	91.31	20	229	0.25	4.23	0.896
5	Tera-10	42.9							
6	Thunderbird	38.27							
7	Fire x4600	38.18							
8	BlueGene eServer	37.33							
9	Red Storm	36.19	32.99	1813.06	43.58	1118	1.02	7.97	1.149
10	Earth Simulator	35.86							40

#### **Balance: Random Ring B/W and CPU Speed**

