

# Performance Tuning and OpenMP

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OpenMP Performance Tuning  
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## Outline

- Motivation
- Performance Basics
- General Performance Issues and OpenMP
- Special Performance Hints for OpenMP

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## Motivation

Reasons for parallel programming:

1. Higher Performance
  - Solve the same problem in shorter time
  - Solve larger problems in the same time
2. Higher Capability
  - Solve problems that cannot be solved on a single processor
    - Larger memory on parallel computers, e.g. 128 GB on hwwsr8k
    - Time constraints limits the possible problem size  
( Weather forecast, turn around within working day)

In both cases performance is one of the major concerns.

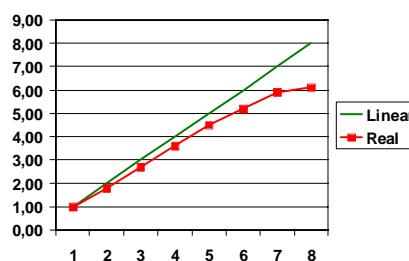
## Performance Basics: Speed Up

- Definition of speed up  $S$

$$S = \frac{T_s}{T_p}$$

$T_s$ : Serial Execution Time       $T_p$ : Parallel Execution Time

- Speed up versus number of used processors:



## Performance Basics: Amdahl's Law

- Assumption:
  - Only a fraction  $F$  of the algorithm is parallel with speed up  $S_p$
  - A fraction  $(1-F)$  is serial
- Total speed up:

$$S = \frac{1}{(1-F) + \frac{F}{S_p}}$$

- Even with infinite parallel speed up your total speed up is limited to:

$$S = \frac{1}{(1-F)}$$

## Consequence of Amdahl's law: necessary parallelization

- If you know your desired speed up  $S$  you can calculate  $F$ :

$$F = 1 - \frac{1}{S}$$

- $F$  gives you the percentage of your program that has to be executed parallel in order to achieve a speed up  $S$
- In order estimate the resulting effort you need to know in which parts of your program  $100*(1-F)\%$  of the time is spent.

## Performance measurement: Profiling

- On most platforms:  
compile with option -p to get program counter profiling

#calls	Time (%)	Accumulated Time (%)	Call
155648	31.22	31.22	Calc
603648	22.24	53.46	Multiply
155648	10.05	63.51	Matmul
214528	9.33	72.84	Copy
603648	7.87	80.71	Find
- Examples:
  - for a speed up of 2 you need to parallelize Calc and Multiply.
  - for a speed up of 5 you need to parallelize Calc, Multiply, Matmul, Copy and Find
- Advantage of OpenMP: this incremental approach is possible

## General issues: Problem size dependency of performance

- Example: Norm of a Matrix:
$$\|A\| = \max_j \sum_i |A_{ij}|$$
- Simple Algorithm:

```
do j=1,n
    b(j)=0
    do i=1,n
        b(j) = b(j) + abs(a(i,j))
    end do
    end do
    do j=1,n
        result = max(result,b(j))
    end do
```
- Change Matrix size from 1 to 4096 and check the performance

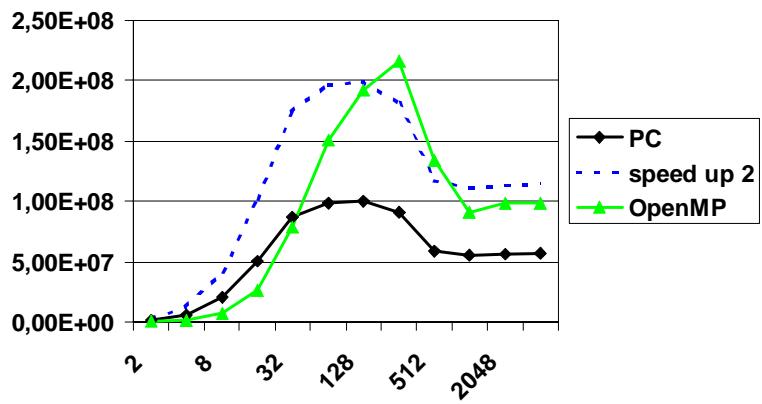
**OpenMP version of Matrix Norm**

```

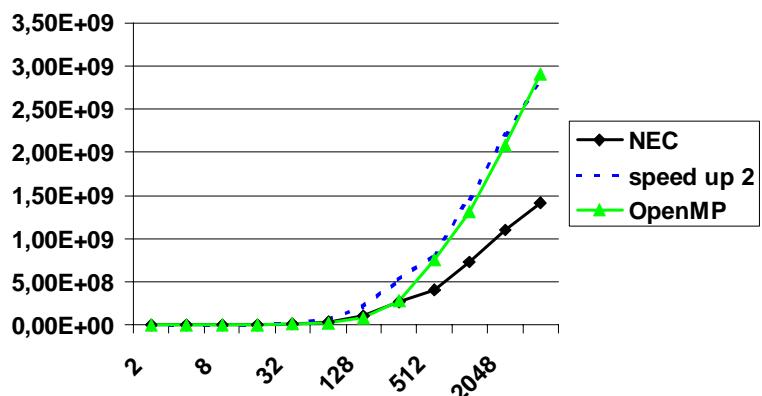
!$OMP PARALLEL
 !$OMP DO PRIVATE(i)
    do j=1,n
        b(j)=0
        do i=1,n
            b(j) = b(j) + abs(a(i,j))
        end do
    end do
 !$OMP DO REDUCTION(MAX:result)
    do j=1,n
        result = max(result,b(j))
    end do
 !$OMP END PARALLEL

```

**Performance on a PC (Dual Pentium II, 450 MHz)**



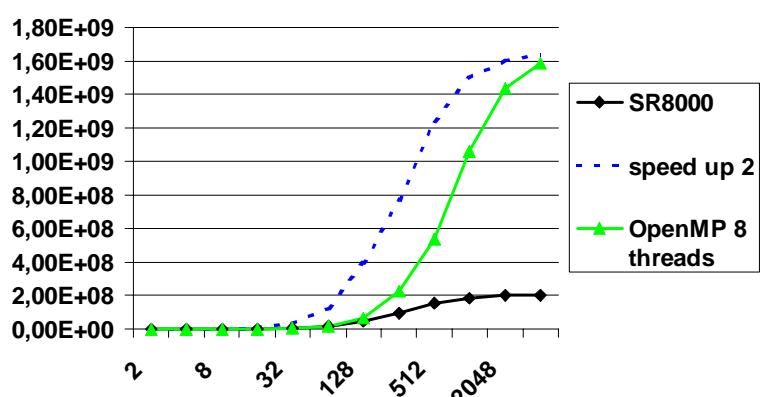
### Performance on a Vectorcomputer (NEC SX-5Be)



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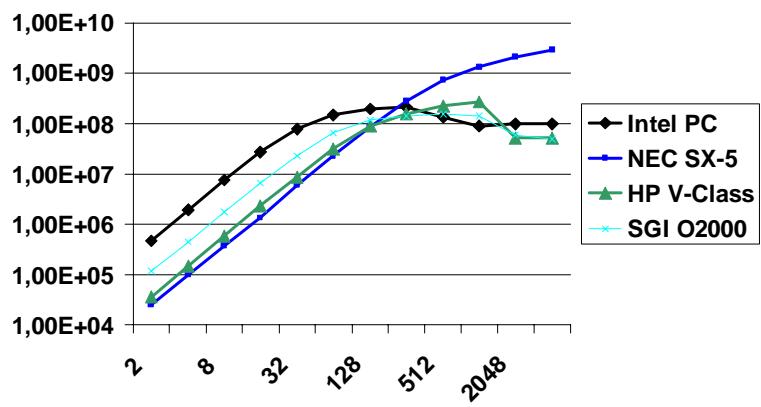
### Performance on the Hitachi SR8000



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### Performance comparison (2 threads with OpenMP)



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### Use OpenMP only with sufficient workload: if-clause

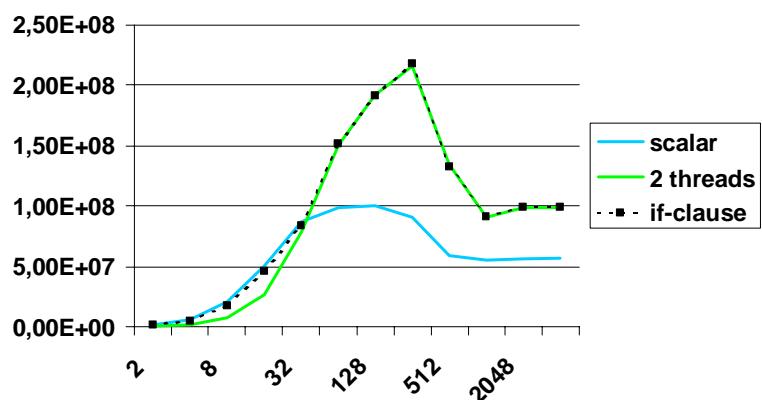
- Only start parallel thread if there is enough workload, otherwise code is executed serial

```
!$OMP PARALLEL IF(n>32)
 !$OMP DO PRIVATE(i)
   do j=1,n
     b(j)=0
     do i=1,n
       b(j) = b(j) + abs(a(i,j))
     end do
   end do
 !$OMP DO REDUCTION(MAX:result)
   do j=1,n
     result = max(result,b(j))
   end do
 !$OMP END PARALLEL
```

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### Performance with if-clause

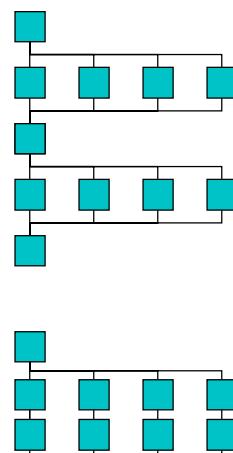


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### Performance with OpenMP: Avoid thread creation

```
#pragma omp parallel for
for(i=0; i<size; i++)
    a[i] = 1.0/a[i];
#pragma omp parallel for
for(i=0; i<size; i++)
    b[i] = b[i]*2.0
```



The improved version only creates the threads once:

```
#pragma omp parallel
{
    #pragma omp for
    for(i=0; i<size; i++)
        a[i] = 1.0/a[i];
    #pragma omp for
    for(i=0; i<size; i++)
        b[i] = b[i]*2.0
}
```

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## Performance with OpenMP: Avoid barriers

### 1. Merge loops:

Replace:

```
#pragma omp for
for(i=0; i<size; i++)
    a[i] = 1.0/a[i];
#pragma omp for
for(i=0; i<size; i++)
    b[i] = b[i]*2.0
```

with

```
#pragma omp for
for(i=0; i<size; i++){
    a[i] = 1.0/a[i];
    b[i] = b[i]*2.0;
}
```

## Performance with OpenMP: Avoid barriers (II)

### 2. Use the NOWAIT clause

- An implicit barrier is put at the end of each work-sharing construct
- It can be eliminated by using nowait, if the barrier is not required

```
#pragma omp parallel
{
#pragma omp for nowait
    for(i=0; i<size; i++)
        a[i] = 1.0/a[i];
#pragma omp for
    for(i=0; i<n; i++)
        b[i] = b[i]*2.0
}
```

## Performance with OpenMP: load balancing

- Different scheduling mechanisms help to avoid load imbalance
- Be aware that dynamic scheduling has a larger overhead
- The default scheduling is implementation dependant, but probably very similar to SCHEDULE( STATIC ), with one chunk for each thread
- Use SCHEDULE( DYNAMIC [ ,*chunksize* ] ) if the workload for each iteration is large and the workload is not predictable
- Use SCHEDULE( STATIC, *chunksize* ) if the load balance can be achieved by reducing the chunksize
- Use SCHEDULE( GUIDED [ ,*chunksize* ] ) as a compromise between STATIC and DYNAMIC
- Use SCHEDULE( RUNTIME ) if the best scheduling depends strongly on the input data, set OMP\_SCHEDULE accordingly

## Summary

- Do not forget serial performance
  - profiling helps to understand the effort to parallelize your program
  - the performance depends on the problem size, using parallelism reduces the problem size on each thread
- Avoid thread creation and unnecessary synchronization
- Use dynamic scheduling strategies only if you have load imbalance problems
- Try to find the best platform for your problem

### Exercise: Matrix Norm Calculation

- Program calculates the norm of a rectangular matrix and a triangular matrix
- Fortran Version: norm\_f90.f90  
C Version: norm\_c.c  
htrs\_get\_time.c: utility function for time measurement

### Exercise: Matrix Norm Calculation on SR8000

#### Step 1: Compile and run the serial version of your program

##### On crosscompiler platform:

- Fortran:
  - xf90 -c norm\_f90.f90 -OSS -noperll norm\_f90.o
  - xcc -c -O4 -pvec +Op -noperll htrs\_get\_time.c
  - xf90 -o norm\_f90 htrs\_get\_time.o norm\_f90.o
- C:
  - xcc -c -O4 -pvec +Op -noperll norm\_c.c
  - xcc -c -O4 -pvec +Op -noperll htrs\_get\_time.c
  - xcc -o norm\_c norm\_c.o htrs\_get\_time.o

##### On Hitachi SR8000:

- run norm\_c or norm\_f90

## Exercise: Matrix Norm Calculation on SR8000

### Step 2: Profile the program

- Fortran:
  - xf90 -c norm\_f90.f90 -OSS -noparallel -Xfuncmonitor norm\_f90.f90
  - xcc -c -O4 -pvec +Op -noparallel htrs\_get\_time.c
  - xf90 -o norm\_f90 htrs\_get\_time.o norm\_f90.o -lpl -parallel
- C:
  - xcc -c -O4 -pvec +Op -noparallel norm\_c.c -Xfuncmonitor
  - xcc -c -O4 -pvec +Op -noparallel htrs\_get\_time.c
  - xcc -o norm\_c norm\_c.o htrs\_get\_time.o -lpl -parallel

### On Hitachi SR8000:

- run norm\_c or norm\_f90 and check pl\_norm\_XXXX.txt

## Exercise: Matrix Norm Calculation on SR8000

### Step 3: Add OpenMP directives

- Concentrate on the most time consuming parts of your program.  
Think about the speed up you like to achieve and the consequences of Amdahl's law.

### Step 4: Compile and run

- Fortran:
  - xf90 -c -OSS -parallel -omp norm\_f90\_omp.f90 -Xfuncmonitor
  - xcc -c -O4 -pvec +Op -parallel -omp htrs\_get\_time.c
  - xf90 -parallel -omp -o norm\_f90 htrs\_get\_time.o norm\_f90\_omp.o -lpl
  - prun -p single ./norm\_f90
- C:
  - xcc -c -O4 -pvec +Op -parallel -omp norm\_c\_omp.c -Xfuncmonitor
  - xcc -c -O4 -pvec +Op -parallel -omp htrs\_get\_time.c
  - xcc -parallel -omp -o norm\_c norm\_c\_omp.o htrs\_get\_time.o -lpl
  - prun -p single ./norm\_c

## Exercise: Matrix Norm Calculation

### Step 5: Optimize

- Try to get rid of redundant synchronization by adding nowait or merging loops
- Try to improve the performance for small matrices, compare the performance with the serial code
- Think about possible load imbalance

Hint: for your convenience there is a gnuplot script file.

- Save the output of the serial program in “scalar.log”
- Save the output of the OpenMP program in “parallel.log”
- Save the output of the optimized version in “optimized.log”
- Compare the results with “gnuplot makeplots.gnu”

## Gnuplot

Hint: for your convenience there is a gnuplot script file.

Fortran

- Save one result in “result.log”
- Save the output of the optimized version in “solution\_f90.log”
- Compare the results with “gnuplot makeplots\_f90.gnu”

C:

- Save one result in “result.log”
- Save the output of the optimized version in “solution\_c.log”
- Compare the results with “gnuplot makeplots\_c.gnu”