

# Object Oriented Parallel Programming with C++

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## Outline

- Why Object Oriented Programming?
- Some Language Philosophy
- Performance issues
- Case Studies:
  - Short Range Molecular Dynamics
  - CFD
  - Particles in liquid
- Conclusion



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## Topology of a Classical Programming Language

Daten:

Unterprogramme:



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## Topology of a Programming Language with Modules

Daten:

Unterprogramme:



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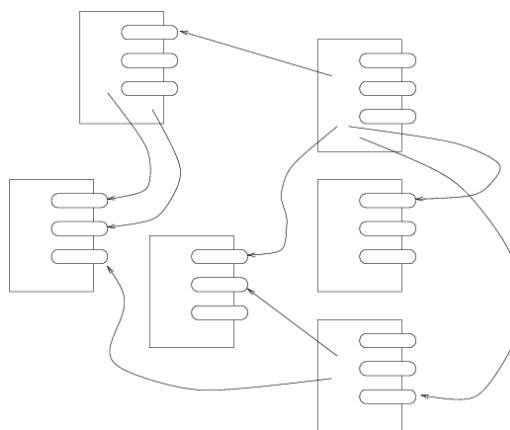
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## Topology of an Object Oriented Language



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## Possible Benefits of Object Oriented Programming

- Better maintainability of programs
- More frequent code re-use
- More efficient software development in groups
- Higher adaptability of software to new demands
- ...

This benefits are especially important for scientific environments with an average student "lifetime" between one and five years.



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## Possible Drawbacks of Object Oriented Programming

- Overhead of interface and object definitions will compensate benefits for small programs.
- Object oriented programming may not be suitable for scientific programming.
- Problem oriented languages might be better than multi-purpose object oriented languages.
- A wrong design will result in a costly re-design of the program instead of a quick hack.
- Abstraction may introduce performance penalty.



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## Language Philosophy

- OO languages to choose from:
  - Smalltalk
  - Eiffel
  - Ada
  - C++
  - Java
  - Fortran 2000?
- Java and C++ are widespread
- Only C++ is available and supported by vendor on Supercomputers like NEC SX, Hitachi SR8000, ...
- It is possible to write Object Oriented programs in “classical” languages, but this requires great self discipline and sometimes results in reduced performance



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## History of C++

- 1985 first commercial release of C++ by AT&T
- 1989 beginning of standardization process
- 1. September 1998 ISO/IEC 14882 language standard.
- The standard is quite new.
- Some language features were added quite late.
- The standard includes the definition of a huge standard library. Which requires very recent language features.
- Some compilers still do not implement a sufficient subset of the standard (SUN CC, IBM xLC).
- Many compilers do not include a complete standard library (Cray CC)
- The last two points may change daily.



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## C++ Language Design

- C++ is a multi-paradigm language
  - Procedural programming:
    - **block structures**
    - **functions**
  - Modular programming:
    - **namespaces**
    - **file scope of identifiers**
  - Object Oriented Programming
    - **classes**
    - **inheritance**
    - **polymorphism**
    - **overloading of functions**
  - Generic Programming
    - **templates**



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## C++ Language Design

- With a few minor exceptions C++ is a better C:  
This allows a smooth migration from C to C++.
- C++ was designed to have optimal run-time efficiency:
  - you don't pay what you don't use
  - polymorphism:
    - C++ uses static polymorphism wherever possible.  
This is necessary to avoid indirect function calls.
    - Run time overhead is under the control of the user.
    - Availability of templates enables programs with flat inheritance hierarchy, or no inheritance at all.
  - templates
    - do not introduce any run-time overhead
    - allows to write generic programs for user and built-in types together with operator overloading
  - no safety checks at run-time



## C++ Language Design

- Memory is under the control of user
  - location of variables (heap, stack) according to lifetime of object
  - user decides whether default initialization is required
  - no garbage collector, de-allocation is explicit in destructor
  - user can provide his own allocation strategy
    - overloading of `new` and `delete`
    - users can provide allocators to containers in STL

The performance of many scientific programs is dominated by memory access!



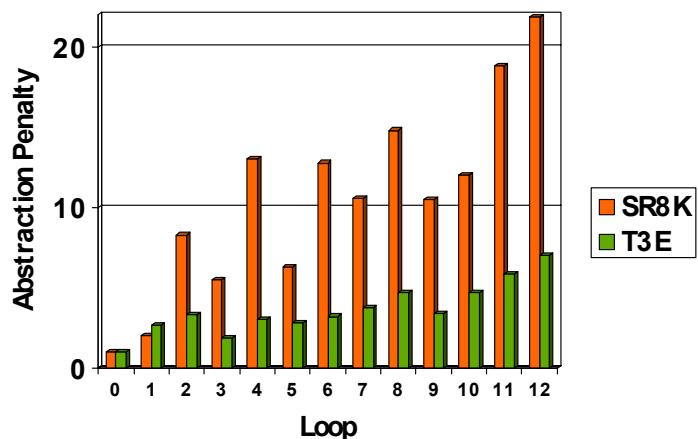
## Performance and C++: Stepanov Abstraction Benchmark

- Well known C++ compiler benchmark (used by KAI and gcc).
- Adds 2000 doubles in an array 25000 times.
- 13 different loops that add more and more abstraction.
- A perfect compiler should generate the same code for all loops.
- Performance is given in MFlops and relative to loop 0 (Fortran style loop)
- Overall compiler quality is given as **abstraction penalty**:  
Geometric mean of performance of the loop 1-12 relative to loop 0.  
**It claims to represent the factor you will be punished by the compiler if you use C++ data abstraction features.**

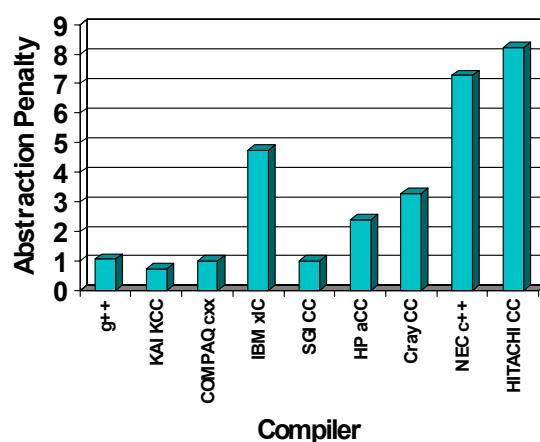
## Stepanov benchmark: description of loops

Nr.	Description
0	Fortran style loop
1-12	STL like accumulators with plus function obj.
1,3,5,7,9,11	doubles
2,4,8,10,12	Doubles wrapped in a class
1,2	Regular pointers
3,4	Pointers wrapped in a class
5,6	Pointers wrapped in a reverse-iterator
7,8	Wrapped Pointers wrapped in adapt.
9,10	Wrapped Pointers double wrapped in adapt.
11,12	Double wrapped Ptrs. Double wrapped in class

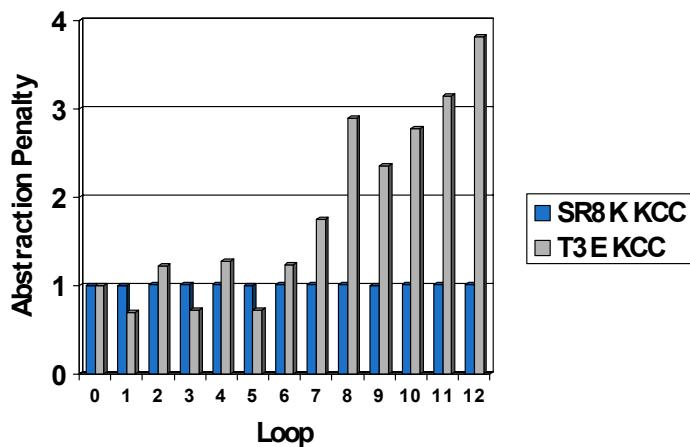
## Results of Stepanov Benchmark on SR8000 and T3E



## Results of Stepanov Benchmark for Different Platforms



## Results of Stepanov Benchmark on SR8000 and T3E



## Availability of KCC at HLRS

- Most Workstation Platforms:
  - Compaq True64, HP HP-UX, IBM AIX, Linux (Redhat, i386), SGI Irix, SUN Solaris
  - Windows NT
- Some Supercomputer Platforms:
  - Cray T3E, Hitachi SR2201, Hitachi SR8000
- KCC is installed on T3E, SR2201 and SR8000 (for T3E “module load KCC”, for others in PATH)
- Floating License for University Stuttgart  
2 licenses are available  
Download the software from <http://www.kai.com>  
(LM\_LICENSE\_FILE=7244@servint1.rus.uni-stuttgart.de)
- Information:  
<http://www.hlrs.de/organization/par/services/tools/compiler/kcc.html>

## C++ Performance Issues

- Multi Language Programming
- Inlining
- Aliasing
- Temporary Objects
- Expression Templates



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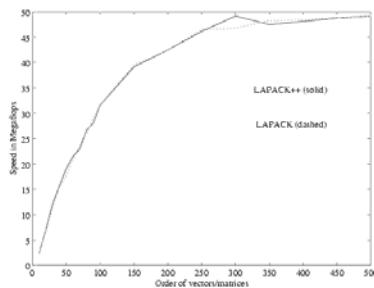
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## C++ and Performance: Multi Language Programming

- Because C++ allows control of memory layout one possible approach is to use C++ for interfaces and complex pre- and post-processing and use Fortran for the numerical expensive part.
- This approach is successfully taken e.g. by Lapack++



- Only suitable if basic data structures are simple, like dense matrices.



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## C++ and Performance: Inlining

- Inlining should eliminate the overhead of function calls
- Classification of functions:
  - Small functions
    - computing time is dominated by function call

```
double& operator[](int i){  
    return data[i];  
}
```
    - Inlining is easy
  - Medium size functions
    - Significant overhead of function call
    - Medium complexity of calculation
    - Inlining is difficult
  - Large size functions
    - Overhead of function call is negligible
    - Large and complex calculations
    - Inlining does not matter, or is even a handicap due to code size



## C++ and Performance: Inlining

- For a highly optimizing compiler there should be no “medium sized functions”.
- But: reasons that can inhibit inlining
  - no source code available: provide definition in header file
  - local variables:

```
write  
    return data[i];
```

instead of

```
double rvalue=data[i];  
return rvalue;
```
  - function has block scopes (loop, if-then-block):  
blame your compiler writer
- some compilers only inline if inline specifier is provided by user



## C++ and Performance: Inlining

- Virtual functions may inhibit inlining, because they are implemented with an indirect function call via the virtual function table

```
class A{  
    virtual f();  
};  
A myA;  
A* p=&myA;  
myA.f(); //might be inlined  
p->f(); // will not be inlined
```

- Use static polymorphism (non virtual functions, templates) instead of dynamic polymorphism wherever possible
- If you really need virtual functions, they will probably be as fast or faster as any hand written code



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## C++ and Performance: Aliasing

- C++ inherited the alias problems from C
- Example

```
void rank1update(Matrix& A, Vector& x){  
    for(int i=0; i<A.rows(); i++)  
        for(int j=0; j<A.cols(); j++)  
            A(i,j)=x(i)*x(j);  
}
```

- Compiler does not know that x is not part of A
- `restrict` qualifier may help, but is not part of C++ standard.  
Some compilers (KCC, Cray, ..) support this keyword.
- with this qualifier the user gives the promise that there is no alias
- in Fortran this promise is part of the language definition
- you may use a macro to eliminate `restrict`



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## C++ and Performance: Problem of Temporary Objects

- C++ allows to write expressive code like:

```
TinyVec v1,v2,v3;  
v1=v1+v2+v3;
```

- A classical code will look like:

```
class TinyVec{  
public:  
    TinyVector operator+(const TinyVec& rhs){  
        TinyVec rvalue=*this;  
        rvalue+=rhs;  
        return rvalue;  
    }  
    const TinyVec& operator=(...);  
    const TinyVec& operator+=(...);  
};
```



## C++ and Performance: Problem of Temporary Objects

- This will generate code similar to:

```
TinyVec v1,v2,v3;  
TinyVec tmp1,tmp2;  
tmp1=operator_plus(v1,v2);  
tmp2=operator_plus(tmp1,v3);  
v1=tmp2;
```

- Possible solutions

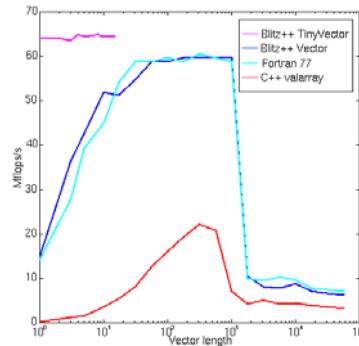
- good compilers will inline all calls and eliminate temporary objects.
- careful introduction of special return types and overloaded functions will eliminate the need of temporaries (complicated)
- Advanced techniques like template expressions will in addition perform loop unrolling etc. (complicated)
- Use a library that implements above optimization techniques.

- First or last approach is recommended



## C++ and Performance: Expression Templates

- Idea: use template techniques to make calculations and loop unrolling at compile time.
- Example: **blitz++** library <http://www.oonumerics.org/blitz>
- Performance of daxpy operation:  $y=y+a*x$



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## C++ and MPI: Template Problem

- Problem: in generic template code you write code for type T. The type can represent anything from double to some user-defined type. As soon as you use external libraries you need to write something like MPI\_DOUBLE if T is a double, and maybe my\_particle\_mpi\_type if T is a user type.
- This problem is not limited to MPI but occurs in all situations, where you need to map a type to some other information (string, upper limit of this type, etc...)
- Solution: traits
- This solution is also used in `numerical_limits` provided in header `<limits>`.

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## C++ and MPI: Generic programming with traits

- Provide template class:  

```
template<class T>
class PTM_MPI_traits {
public:
    static inline MPI_Datatype datatype(void);
};
```
- Specialize for every type you need:  

```
template<>
class PTM_MPI_traits<double>{
public:
    static inline MPI_Datatype datatype(void){return MPI_DOUBLE;};
};
```
- Use the traits in MPI Calls:  

```
MPI_Send(...,PTM_MPI_traits<T>::datatype(),...);
```



## Case Studies

- Molecular Dynamics
- CFD Code
- Sedimentation of Particles in Liquid



## Molecular Dynamics

- Different particle types
  - spheres
  - ellipsoids
  - polygons , ....
- Different forces
  - Lennard-Jones
  - Contact Forces
- Different Integrators
  - Verlet, velocity Verlet
  - Leap Frog
  - Predictor Corrector
- Versatile data analysis



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## Molecular Dynamics

### Particles

- Location
- Velocity
- Mass

### Container

- Size
- Boundary Conditions
- Particles → Linked Cell

### Functions

- Force Calculation
- Integration
- Analysis

### Algorithms

- for each particle
- for each pair



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## Main loop

```
double dt=0.1; // time step
PartContLC<Particle,3> box(ll,ur,bound_cond,cut_off);
Int_vv_start<Particle> int_start(dt); // integration
Int_vv_finish<Particle> int_finish(dt);

while( t < maxT ){
    // velocity verlet integration fist part
    for_each_particle(box.begin(),box.end(),int_start);
    // force calculation
    box.update();
    box.for_each_pair(myForce);
    // velocity verlet integration second part
    for_each_particle(box.begin(),box.end(),int_finish);
    t+=dt;
}
```

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## Forces

```
template<class PARTICLE>
class Force{
public:
    inline void operator()(PARTICLE& p1, PARTICLE& p2){
        Vector force;
        // calculation of f goes here
        p1.f += force;
        p2.f -= force; // actio = reactio
    }
};

Force<myParticle> myForce;
myParticle p1,p2;
myForce(p1,p2);
```

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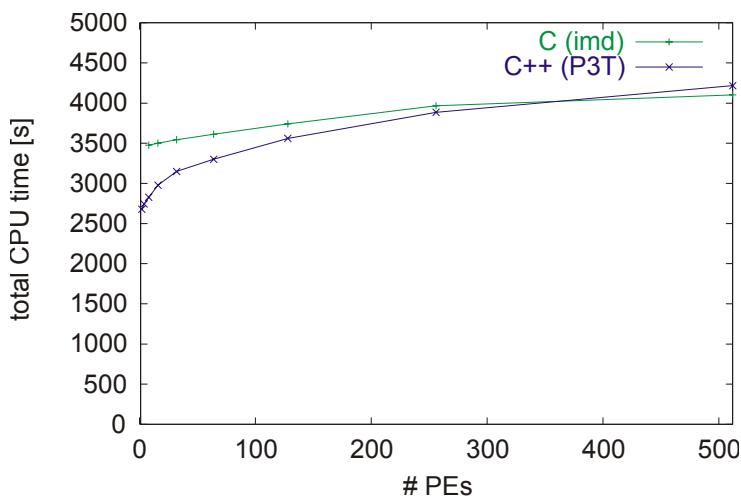
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## Performance in Comparison to C



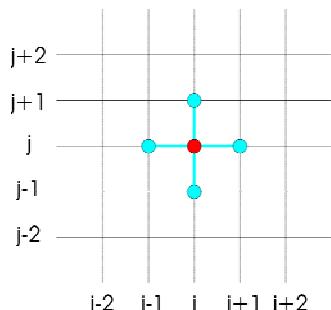
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## Partial Differential Equations

### Partial Diff. Eq.:

- local information
- rectangular grid
- n-dimensional



Example: Poisson-Equation

$$\nabla^2 p = f(\mathbf{x})$$

discretized:

$$\frac{1}{h^2} \left( p_{i-1,j,k} + p_{i+1,j,k} + p_{i,j-1,k} + p_{i,j+1,k} + p_{i,j,k-1} + p_{i,j,k+1} - 6p_{i,j,k} \right) = f_{i,j,k}$$

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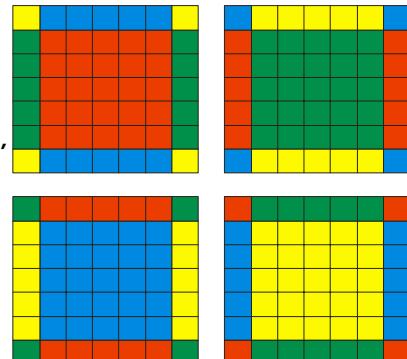
## n-dimensional parallel Grid

```
const int dimension=2;
const IntVector cpu_layout(2,2);
const int n_shadow=1;
const intVector boundary_cd(1,1);
const intVector size(10,10);

// define PE layout
PT_Parallel PEs(dimension,cpu_layout,
                  boundary_cd);

// create field
PT_PArray<double> p(PEs, size,
                      n_shadows);

// initialize values
init_field(p);
// update boundary conditions
p.update_bc();
```



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## Gauss-Seidel Smoother

```
struct Poisson_Node { double p, double f };

PT_PArray<Poisson_Node> p( .... );
GS_operator<Poisson_Node> GaussSeidel(h);

while ( error > limit ){
    // apply boundary conditions
    p.update_bc();
    // apply Gauss Seidel operator
    // alternatingly linewise
    apply_aw(p,p.begin(),p.end(),GaussSeidel);
}
```



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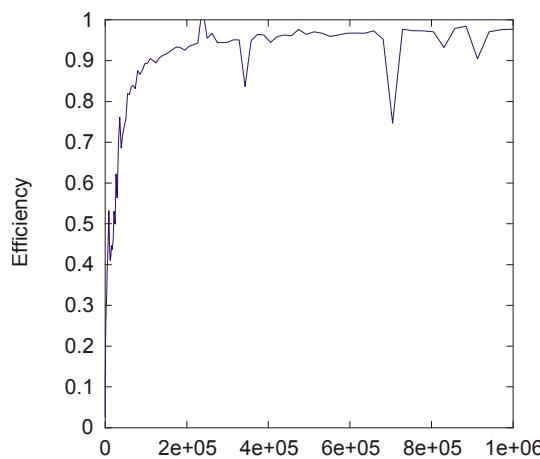
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## Parallel Efficiency

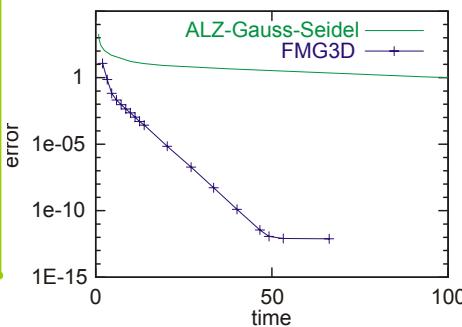


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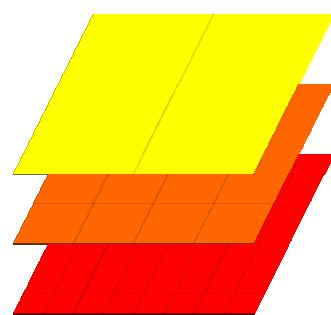
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## Multigrid Solver

The use of a multigrid solver improves the poor scaling behavior of the iterative Gauss Seidel method.



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## Multigrid Solver

```
const IntVector    cpu_layout(2,4,2);
const IntVector    boundary_cd(1,0,1);
const int          dimension=3;
const int          n_shadow=1;
const IntVector    field_size(257,513,257);

PT_Parallel       PE_layout(Dimension,cpu_layout,boundary_cd);
PT_PArray<double> pressure(PE_layout,field_size,n_shadows);
PT_PArray<double> div_f(PE_layout,field_size,n_shadows);
PT_Poisson        poisson(PE_layout,field_size,n_shadows);

// ... compute rhs (div_f) and solve Poisson equation
poisson.rhs() = div_f;
poisson.solve();
pressure      = poisson.x();
ofstream      save("restart.dat");
save << pressure;
```



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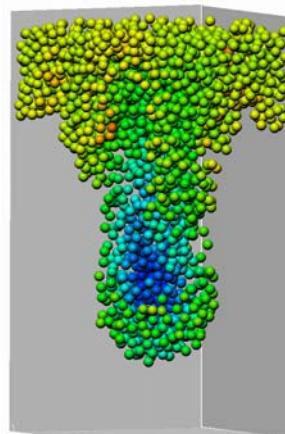
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## Sedimentation of Particles in Liquid

- Relevant system size:  $10^5$ - $10^8$  particles
- Numerical treatment is expensive
- Treatment of subproblems
  - Molecular Dynamics
    - Forces
  - Partial Differential Equations
    - Navier-Stokes equation with moderate Reynolds numbers
    - parallel / serial data structures
    - Multigrid Solver for Poisson Equation
    - moving boundaries  
(explicit force density)



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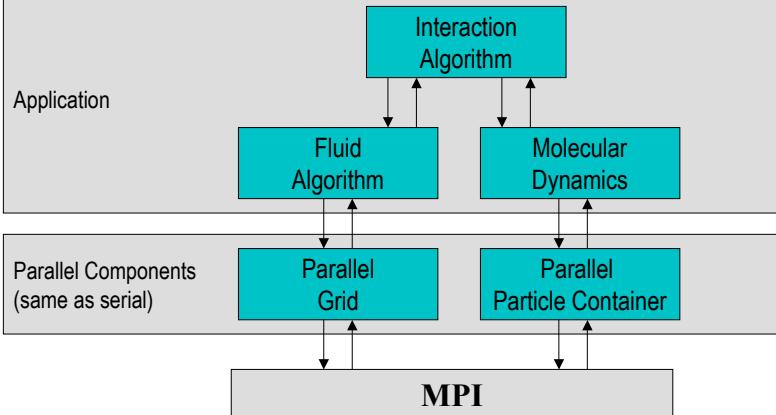
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## Structure of Particles in Liquid Simulation Program



## Conclusion

- Object Oriented programming in C++ is possible without performance penalty
- Separation of algorithms and data results in greater flexibility
- Encapsulation of parallelism should be one goal of design